# 68'

## MICRO JOURNAL

## OS-9 Atari Amiga Mac S-50

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**OS-9**  SK*DOS  Atari Amiga
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## VOLUME X   ISSUE VIII ● Devoted to the 68XXX User ● August 1988

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## 68 MICRO JOURNAL

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# C

## The C Programmers Reference Source. Always Right On Target!

## C User Notes

## A Tutorial Series

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### INTRODUCTION

This chapter continues the discussion of dbug, a C debugging package. It is a useful tool for debugging and testing C programs. It was developed by Fred Fish, who placed it into the public domain. The C code for the dbug package and for extensions to it appear in this and subsequent chapters.

### DBUG.H

The Dbug.h header file must be included in any program in case Fred Fish's dbug package is to be used. It contains the appropriate preprocessor commands to call support routines in the dbug run-time library contained in dbug.c.

To disable compilation of the dbug commands, define the preprocessor symbol "DBUG_OFF". This will result in null expansions so that the resulting code will be smaller.

All externally visible symbol names follow the pattern _db_xxx..xx_ to minimize the possibility of a dbug package symbol coinciding with a user-defined symbol.

```
/*
 *      dbug.h      header file for dbug package
 *
 *      Fred Fish
 */

/*
 *      Internally used variables which must be global.
 */

#ifndef DBUG_OFF
extern int   _db_on_;          /* TRUE if debug currently enabled */
extern FILE *_db_fp_;          /* Current debug output stream */
extern char *_db_process_;     /* Name of current process */
extern int   _db_keyword_();   /* Accept/reject keyword */
extern void  _db_push_();      /* Push state, set up new state */
extern void  _db_pop_();       /* Pop previous debug state */
extern void  _db_enter_();     /* New user function entered */
extern void  _db_return_();    /* User function return */
extern void  _db_pargs_();     /* Remember args for line */
extern void  _db_doprnt_();    /* Print debug output */
extern void  _db_setjmp_();    /* Save debugger environment */
extern void  _db_longjmp_();   /* Restore debugger environment */
#endif
```

```
/*
 *      These macros provide a user interface into functions in the
 *      dbug runtime support library.  They isolate users from changes
 *      in the MACROS and/or runtime support.
 *
 *      The symbols "__LINE__" and "__FILE__" are expanded by the
 *      preprocessor to the current source file line number and file
 *      name respectively.
 *
 *      WARNING - Because the DBUG_ENTER macro allocates space on
 *      the user function's stack, it must follow any declarations and
 *      precede any executable statements in the user function.
 *
 */

#ifdef DBUG_OFF
#define DBUG_ENTER(a1)
#define DBUG_EXECUTE(keyword,a1)
#define DBUG_FILE(stderr)
#define DBUG_LONGJMP longjmp
#define DBUG_POP()
#define DBUG_PRINT(keyword,arglist)
#define DBUG_PROCESS(a1)
#define DBUG_PUSH(a1)
#define DBUG_RETURN(a1) return(a1)
#define DBUG_SETJMP setjmp
#define DBUG_VOID_RETURN return
#else
#define DBUG_ENTER(a) char *_db_func_, *_db_file_; int _db_level_; \
        _db_enter_(a,__FILE__,__LINE__,&_db_func_,&_db_file_,&_db_level_)
#define DBUG_EXECUTE(keyword,a1) \
        { if (_db_on_) { if (_db_keyword_(keyword)) { a1 } } }
#define DBUG_FILE (_db_fp_)
#define DBUG_LEAVE (_db_return_(__LINE__, &_db_func_, &_db_file_,
&_db_level_))
#define DBUG_LONGJMP(a1,a2) (_db_longjmp_(), longjmp(a1, a2))
#define DBUG_POP() _db_pop_()
#define DBUG_PRINT(keyword,arglist) \
        { if (_db_on_) { _db_pargs_(__LINE__,keyword); _db_doprnt_ ar-
glist; } }
#define DBUG_PROCESS(a1) (_db_process_ = a1)
#define DBUG_PUSH(a1) _db_push_(a1)
#define DBUG_RETURN(a1) {DBUG_LEAVE; return(a1);}
#define DBUG_SETJMP(a1) (_db_setjmp_(), setjmp(a1))
#define DBUG_VOID_RETURN {DBUG_LEAVE; return;}
#endif
```

DBUG.C

The code below represents the runtime library supporting the dbug
package.

```
/*
 *      dbug.c    runtime support routines for dbug package
 *
 *      These are the runtime support routines for the dbug package.
 *      The dbug package has two main components; the user include
 *      file containing various macro definitions, and the runtime
 *      support routines which are called from the macro expansions.
 *
 *      Externally visible functions in the runtime support module
 *      use the naming convention pattern "_db_xx...xx_", thus
 *      they are unlikely to collide with user defined function names.
 *
 *      Fred Fish                (base code)
 *      (Currently at Motorola Computer Division, Tempe, Az.)
 *      hao!noao!mcdsun!fnf
 *      (602) 438-3614
 *
 *      Binayak Banerjee         (profiling enhancements)
 *      seismo!bpa!sjuvax!bbanerje
 */

#include <stdio.h>

#ifdef AMIGA
#define H2 (50)                  /* Probably in some header somewhere
*/
#endif

/*
 *      Manifest constants that should not require any changes.
 */

#define FALSE           0        /* Boolean FALSE */
#define TRUE            1        /* Boolean TRUE */
#define EOS             0        /* End Of String marker */

/*
 *      Manifest constants which may be "tuned" if desired.
 */

#define PRINTBUF        1024     /* Print buffer size */
#define INDENT          4        /* Indentation per trace level */
#define MAXDEPTH        200      /* Maximum trace depth default */

/*
 *      The following flags are used to determine which
 *      capabilities the user has enabled with the state push macro.
 */
```

```
    #define DEBUG_ON        000002  /* Debug enabled */
    #define DEPTH_ON        000020  /* Function nest level print enabled
*/
    #define FILE_ON         000004  /* File name print enabled */
    #define LINE_ON         000010  /* Line number print enabled */
    #define NUMBER_ON       000100  /* Number each line of output */
    #define PROCESS_ON      000040  /* Process name print enabled */
    #define PROFILE_ON      000200  /* Print out profiling code */
    #define TRACE_ON        000001  /* Trace enabled */

    #define DEBUGGING (stack -> flags & DEBUG_ON)
    #define PROFILING (stack -> flags & PROFILE_ON)
    #define STREQ(a,b) (strcmp(a,b) == 0)
    #define TRACING (stack -> flags & TRACE_ON)

    /*
    *       Make it easy to change storage classes if necessary.
    */

    #define AUTO auto               /* Names to be allocated on stack */
    #define BOOLEAN int             /* True or false */
    #define EXPORT                  /* Allocated here, available globally
*/
    #define IMPORT extern           /* Names defined externally */
    #define LOCAL static            /* Names not needed by outside world
*/
    #define REGISTER register       /* Names to be placed in registers */
    #define VOID void               /* Some classes have no class */

    /*
    *       The following define is for the variable arguments kluge.
    *       See the comments in _db_doprnt_().
    *
    *       Also note that the longer this list, the less prone to failing
    *       on long argument lists, but the more stuff that must be moved
    *       around for each call to the runtime support routines.  The
    *       length may really be critical if the machine convention is
    *       to pass arguments in registers.
    *
    *       Note that the default define allows up to 16 integral argu-
ments,
    *       or 8 floating point arguments (doubles), on most machines.
    *
    */

    #define ARGLIST a0,a1,a2,a3,a4,a5,a6,a7,a8,a9,a10,a11,a12,a13,a14,a15

    /*
    * The default file for profiling.
    */

    #define PROF_FILE       "dbugmon.out"
```

```
/*
*       Variables which are available externally but should only
*       be accessed via the macro package facilities.
*/

    EXPORT BOOLEAN _db_on_ = FALSE;         /* TRUE if debugging currently
on */
    EXPORT BOOLEAN _db_pon_ = FALSE;        /* TRUE if debugging currently
on */
    EXPORT FILE *_db_fp_ = stderr;          /* Output stream, default
stderr */
    EXPORT FILE *_db_pfp_ = (FILE *)0;      /* Profile stream,
'dbugmon.out' */
    EXPORT char *_db_process_ = "dbug";     /* Pointer to process name;
argv[0] */

/*
*       Externally supplied functions.
*/

#ifdef unix                         /* Only needed for unix */
IMPORT VOID perror ();              /* Print system/library error */
IMPORT int access ();               /* Test file for access */
IMPORT int chown ();                /* Change owner of a file */
IMPORT int getgid ();               /* Get real group id */
IMPORT int getuid ();               /* Get real user id */
#else
#if !(AMIGA && LATTICE)
LOCAL VOID perror ();               /* Fake system/library error print
routine */
#endif
#endif

#if BSD4_3 || sun
IMPORT int getrusage ();
#endif

IMPORT FILE *fopen ();              /* Open a stream */
IMPORT VOID exit ();                /* Terminate execution */
IMPORT VOID free ();
IMPORT char *malloc ();             /* Allocate memory */
IMPORT char *strcpy ();             /* Copy strings around */
IMPORT int atoi ();                 /* Convert ascii to integer */
IMPORT int fclose ();               /* Close a stream */
IMPORT int fprintf ();              /* Formatted print on file */
IMPORT int strcmp ();               /* Compare strings */
IMPORT int strlen ();               /* Find length of string */

#ifndef fflush                      /* This is sometimes a macro */
IMPORT int fflush ();               /* Flush output for stream */
#endif
```

```
    /*
     *      The user may specify a list of functions to trace or
     *      debug.  These lists are kept in a linear linked list,
     *      a very simple implementation.
     */

    struct link
    {
        char *string;                   /* Pointer to link's contents */
        struct link *next_link;         /* Pointer to the next link */
    };


    /*
     *      Debugging states can be pushed or popped off of a
     *      stack which is implemented as a linked list.  Note
     *      that the head of the list is the current state and the
     *      stack is pushed by adding a new state to the head of the
     *      list or popped by removing the first link.
     */

    struct state
    {
        FILE *out_file;                     /* Current output stream */
        FILE *prof_file;                    /* Current profiling stream */
        int flags;                          /* Current state flags */
        int level;                          /* Current function nesting
level */
        int maxdepth;                       /* Current maximum trace depth
*/
        struct link *functions;             /* List of functions */
        struct link *keywords;              /* List of debug keywords */
        struct link *p_functions;           /* List of profiled functions
*/
        struct link *processes;             /* List of process names */
        struct state *next_state;           /* Next state in the list */
        unsigned int delay;                 /* Delay after each output
line */
    };

    LOCAL struct state *stack = NULL;       /* Linked list of stacked
states */

    /*
     *      Local variables not seen by user.
     */

    LOCAL BOOLEAN init_done = FALSE;/* Set to TRUE when initialization
done */
    LOCAL char *file = "?file";     /* Name of current user file */
    LOCAL char *func = "?func";     /* Name of current user function */
    LOCAL int lineno = 0;           /* Current debugger output line number
*/

    #if unix || AMIGA
```

```
    LOCAL char *jmpfile;            /* Remember current file for setjmp */
    LOCAL char *jmpfunc;           /* Remember current function for
setjmp */
    LOCAL int jmplevel;            /* Remember nesting level at setjmp ()
*/

    #endif

    LOCAL BOOLEAN DoTrace ();       /* Test for tracing enabled */
    LOCAL BOOLEAN Writable ();      /* Test to see if file is writable */
    LOCAL VOID ChangeOwner ();      /* Change file owner and group */
    LOCAL VOID CloseFile ();        /* Close debug output stream */
    LOCAL VOID DoPrefix ();         /* Print debugger line prefix */
    LOCAL VOID FreeList ();         /* Free memory from linked list */
    LOCAL VOID Indent ();           /* Indent line to specified indent */
    LOCAL VOID OpenFile ();         /* Open debug output stream */
    LOCAL VOID OpenProfile ();      /* Open profile output stream */
    LOCAL VOID PushState ();        /* Push current debug state */
    LOCAL char *BaseName ();        /* Remove leading pathname components
*/

    LOCAL char *DbugMalloc ();      /* Allocate memory for runtime support
*/

    LOCAL char *StrDup ();          /* Make a fresh copy of a string */
    LOCAL struct link *ListParse ();/* Parse a debug command string */
    LOCAL unsigned long Clock ();   /* Return current user time (ms) */

    /* Supplied in Sys V runtime environ */
    LOCAL char *strrchr ();         /* Find last occurance of char */
    LOCAL char *strtok ();          /* Break string into tokens */

    /*
     *      The following local variables are used to hold the state
information
     *      between the call to _db_pargs_() and _db_doprnt_(), during
     *      expansion of the DBUG_PRINT macro.  This is the only macro
     *      that currently uses these variables.  The DBUG_PRINT macro
     *      and the new _db_doprnt_() routine replace the older DBUG_N
macros
     *      and their corresponding runtime support routine _db_printf_().
     *
     *      These variables are currently used only by _db_pargs_() and
     *      _db_doprnt_().
     */

    LOCAL char *u_keyword = "?";    /* Keyword for current macro */
    LOCAL int u_line = 0;           /* User source code line number */

    /*
     *      Miscellaneous printf format strings.
     */

    #define ERR_ABORT "%s: debugger aborting because %s\n"
    #define ERR_CHOWN "%s: can't change owner/group of \"%s\": "
    #define ERR_CLOSE "%s: can't close debug file: "
    #define ERR_MISSING_RETURN "%s: missing DBUG_RETURN macro in function
```

```
\"%s\"\n"
    #define ERR_OPEN "%s: can't open debug output stream \"%s\": "
    #define ERR_PRINTF "%s: obsolete object file for '%s', please
recompile!\n"

    /*
     *      Macros and defines for testing file accessibility under UNIX.
     */

    #ifdef unix
    #define A_EXECUTE       01              /* Test for execute permission
*/
    #define A_EXISTS        00              /* Test for file existance */
    #define A_READ          03              /* Test for read access */
    #define A_WRITE         02              /* Test for write access */
    #define EXISTS(pathname) (access (pathname, A_EXISTS) == 0)
    #define WRITABLE(pathname) (access (pathname, A_WRITE) == 0)
    #else
    #define EXISTS(pathname) (FALSE)        /* Assume no existance */
    #endif

    /*
     *      Translate some calls among different systems.
     */

    #ifdef unix
    #define Delay sleep
    IMPORT unsigned int sleep ();   /* Pause for given number of seconds
*/
    #endif

    #ifdef AMIGA
    IMPORT int Delay ();            /* Pause for given number of ticks */
    #endif

    /*
     *      _db_push_       push current debugger state and set up new one
     *
     *      VOID _db_push_ (control)
     *      char *control;
     *
     *      Given pointer to a debug control string in "control", pushes
     *      the current debug state, parses the control string, and sets
     *      up a new debug state.
     *
     *      The only attribute of the new state inherited from the previ-
ous
     *      state is the current function nesting level.  This can be
     *      overridden by using the "r" flag in the control string.
     *
     *      The debug control string is a sequence of colon separated
fields
     *      as follows:
     *
     *              <field_1>:<field_2>:...:<field_N>
     *
     *      Each field consists of a mandatory flag character followed by
     *      an optional "," and comma separated list of modifiers:
     *
     *              flag[,modifier,modifier,...,modifier]
     *
     *      The currently recognized flag characters are:
     *
     *              d       Enable output from DBUG_<N> macros for
     *                      for the current state.  May be followed
     *                      by a list of keywords which selects output
     *                      only for the DBUG macros with that keyword.
     *                      A null list of keywords implies output for
     *                      all macros.
     *
     *              D       Delay after each debugger output line.
     *                      The argument is the number of tenths of
seconds
     *                      to delay, subject to machine capabilities.
     *                      I.E.  -#D,20 is delay two seconds.
     *
     *              f       Limit debugging and/or tracing, and profiling
to the
     *                      list of named functions.  Note that a null
list will
     *                      disable all functions.  The appropriate "d" or
"t"
     *                      flags must still be given, this flag only
limits their
     *                      actions if they are enabled.
     *
     *              F       Identify the source file name for each
     *                      line of debug or trace output.
     *
     *              g       Enable profiling.  Create a file called
'dbugmon.out'
     *                      containing information that can be used to
profile
     *                      the program.  May be followed by a list of
keywords
     *                      that select profiling only for the functions
in that
     *                      list.  A null list implies that all functions
are
     *                      considered.
     *
     *              L       Identify the source file line number for
     *                      each line of debug or trace output.
     *
     *              n       Print the current function nesting depth for
     *                      each line of debug or trace output.
```

```
*              N         Number each line of debug output.
*
*              P         Limit debugger actions to specified processes.
*                        A process must be identified with the
*                        DBUG_PROCESS macro and match one in the list
*                        for debugger actions to occur.
*
*              P         Print the current process name for each
*                        line of debug or trace output.
*
*              r         When pushing a new state, do not inherit
*                        the previous state's function nesting level.
*                        Useful when the output is to start at the
*                        left margin.
*
*              t         Enable function call/exit trace lines.
*                        May be followed by a list (containing only
*                        one modifier) giving a numeric maximum
*                        trace level, beyond which no output will
*                        occur for either debugging or tracing
*                        macros.  The default is a compile time
*                        option.
*
*      Some examples of debug control strings which might appear
*      on a shell command line (the "-#" is typically used to
*      introduce a control string to an application program) are:
*
*              -#d:t
*              -#d:f,main,subr1:F:L:t,20
*              -#d,input,output,files:n
*
*      For convenience, any leading "-#" is stripped off.
*/


VOID _db_push_ (control)
char *control;
{
    REGISTER char * scan;
    REGISTER struct link * temp;

    if (control && *control == '-')
    {
        if (*++control == '#')
            control++;
    }
    control = StrDup (control);
    PushState ();
    scan = strtok (control, ":");
    for (; scan != NULL; scan = strtok ((char *)NULL, ":"))
    {
        switch (*scan++)
        {
        case 'd':
            _db_on_ = TRUE;
            stack ->flags |= DEBUG_ON;
            if (*scan++ == ',')
                stack ->keywords = ListParse (scan);
            break;
        case 'D':
            stack ->delay = 0;
            if (*scan++ == ',')
            {
                temp = ListParse (scan);
                stack ->delay = DelayArg (atoi (temp ->string));
                FreeList (temp);
            }
            break;
        case 'f':
            if (*scan++ == ',')
                stack ->functions = ListParse (scan);
            break;
        case 'F':
            stack ->flags |= FILE_ON;
            break;
        case 'g':
            _db_pon_ = TRUE;
            OpenProfile(PROF_FILE);
            stack ->flags |= PROFILE_ON;
            if (*scan++ == ',')
                stack ->p_functions = ListParse (scan);
            break;
        case 'L':
            stack ->flags |= LINE_ON;
            break;
        case 'n':
            stack ->flags |= DEPTH_ON;
            break;
        case 'N':
            stack ->flags |= NUMBER_ON;
            break;
        case 'o':
            if (*scan++ == ',')
            {
                temp = ListParse (scan);
                OpenFile (temp ->string);
                FreeList (temp);
            }
            else {
                OpenFile ("-");
            }
            break;
        case 'p':
            if (*scan++ == ',')
                stack ->processes = ListParse (scan);
            break;
        case 'P':
            stack ->flags |= PROCESS_ON;
```

```
                break;
        case 'r':
            stack ->level = 0;
            break;
        case 't':
            stack ->flags |= TRACE_ON;
            if (*scan++ == ',')
            {
                temp = ListParse (scan);
                stack ->maxdepth = atoi (temp ->string);
                FreeList (temp);
            }
            break;
        }
    }
    free (control);
}


/*
 *      _db_pop_    pop the debug stack
 *
 *      Pops the debug stack, returning the debug state to its
 *      condition prior to the most recent _db_push_ invocation.
 *      Note that the pop will fail if it would remove the last
 *      valid state from the stack.  This prevents user errors
 *      in the push/pop sequence from screwing up the debugger.
 *      Maybe there should be some kind of warning printed if the
 *      user tries to pop too many states.
 */

VOID _db_pop_ ()
{
    REGISTER struct state * discard;

    discard = stack;
    if (discard != NULL && discard ->next_state != NULL)
    {
        stack = discard ->next_state;
        _db_fp_ = stack ->out_file;
        _db_pfp_ = stack ->prof_file;
        if (discard ->keywords != NULL)
            FreeList (discard ->keywords);
        if (discard ->functions != NULL)
            FreeList (discard ->functions);
        if (discard ->processes != NULL)
            FreeList (discard ->processes);
        if (discard ->p_functions != NULL)
            FreeList (discard ->p_functions);
        CloseFile (discard ->out_file);
        CloseFile (discard ->prof_file);
        free ((char *) discard);
    }
}
```

```
/*
 *      _db_enter_    process entry point to user function
 *
 *      VOID _db_enter_ (_func_, _file_, _line_, _sfunc_, _sfile_,
 *  _slevel_)
 *      char * _func_;           points to current function name
 *      char * _file_;           points to current file name
 *      int _line_;              called from source line number
 *      char ** _sfunc_;         save previous _func_
 *      char ** _sfile_;         save previous _file_
 *      int * _slevel_;          save previous nesting level
 *
 *      Called at the beginning of each user function to tell
 *      the debugger that a new function has been entered.
 *      Note that the pointers to the previous user function
 *      name and previous user file name are stored on the
 *      caller's stack (this is why the ENTER macro must be
 *      the first "executable" code in a function, since it
 *      allocates these storage locations).  The previous nesting
 *      level is also stored on the callers stack for internal
 *      self consistency checks.
 *
 *      Also prints a trace line if tracing is enabled and
 *      increments the current function nesting depth.
 *
 *      Note that this mechanism allows the debugger to know
 *      what the current user function is at all times, without
 *      maintaining an internal stack for the function names.
 *
 */

VOID _db_enter_ (_func_, _file_, _line_, _sfunc_, _sfile_, _slevel_)
char * _func_;
char * _file_;
int _line_;
char ** _sfunc_;
char ** _sfile_;
int * _slevel_;
{
    if (!init_done)
        _db_push_ ("");
    * _sfunc_ = func;
    * _sfile_ = file;
    func = _func_;
    file = BaseName (_file_);
    stack ->level++;
    * _slevel_ = stack ->level;
    if (DoProfile ())
    {
        (VOID) fprintf (_db_pfp_, "%s\tE\t%ld\n", func, Clock());
        (VOID) fflush (_db_pfp_);
    }
```

```
        if (DoTrace ())
        {
            DoPrefix (_line_);
            Indent (stack ->level);
            (VOID) fprintf (_db_fp_, ">%s\n", func);
            (VOID) fflush (_db_fp_);
            (VOID) Delay (stack ->delay);
        }
    }


/*
 *      _db_return_     process exit from user function
 *
 *      VOID _db_return_ (_line_, _sfunc_, _sfile_, _slevel_)
 *      int _line_;             current source line number
 *      char **_sfunc_;         where previous _func_ is to be retrieved
 *      char **_sfile_;         where previous _file_ is to be retrieved
 *      int *_slevel_;          where previous level was stashed
 *
 *      Called just before user function executes an explicit or implicit
 *      return.  Prints a trace line if trace is enabled, decrements
 *      the current nesting level, and restores the current function and
 *      file names from the defunct function's stack.
 */
VOID _db_return_ (_line_, _sfunc_, _sfile_, _slevel_)
int _line_;
char **_sfunc_;
char **_sfile_;
int *_slevel_;
{
    if (!init_done)
        _db_push_ ("");
    if (stack ->level != *_slevel_ && (TRACING || DEBUGGING || PROFILING))
        (VOID) fprintf (_db_fp_, ERR_MISSING_RETURN, _db_process_, func);
    else if (DoProfile ())
        (VOID) fprintf (_db_pfp_, "%s\tX\t%ld\n", func, Clock());
    else if (DoTrace ())
    {
        DoPrefix (_line_);
        Indent (stack ->level);
        (VOID) fprintf (_db_fp_, "<%s\n", func);
    }
    (VOID) fflush (_db_fp_);
    (VOID) Delay (stack ->delay);
    stack ->level = *_slevel_ - 1;
    func = *_sfunc_;
    file = *_sfile_;
}


The code for the dbug package is continued in the next chapter.

EOF
```

**68 MICRO JOURNAL**™

# Basically OS-9

**A Tutorial Series**

Dedicated to the serious OS-9 user.
The fastest growing users group world-wide!
6809 - 68020

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# DEVICES

I thought this month we could talk about device descriptors. I am sure your saying what more can be said about them. They are a module used by device managers to interact with some piece of hardware. Well there it is in a nutshell. Its been a fun. See you all next time!

*OK. Ok.* There are a few more things that can be said. ( I thought you would never ask. ) Talking about device descriptors is not the most exciting aspect of OS-9, but it is most necessary. So on with it.

I think it is first necessary to understand the nature of peripherals. Those are the things we attach to the computer and expect it to interact with. These include items like monitors, keyboards, terminals, modems, printers, disk drives, and hard disks. The all have one thing in common. They require a communication link with the computer.

OS-9 is clever. It devised a method to talk to all these in a simple and straight forward fashion. Everything boils down to input and output. To handle this OS-9 has a special program called the Input/Output Manger or just IOMan.

In the simplest sense data flows back and forth. Characters flow to your video monitor and from your keyboard. This type of interaction is know as sequential since everything flows one byte after another. To handle this OS-9 has a program called the Sequential Character File Manger or SCFMan for short.

This method works fine for terminals and printers, but it is not practical for disk drives and hard disks. These devices move data in chunks. This chunk is more technically known as a block which consists of a fixed amount of data bytes. The blocks of information are stored in apparent "random" fashion on the device's storage media. ( The word "random" makes me feel that things are stored in a helter-skelter fashion. If this were the case we would not get very far. ) The program that handles this is called the Random Block File Manager or ( you guessed it! ) RBFMan.

There is also a Pipe File Manager called PIPEMan. It handles pipes. Can you guess what the SBFMan is? Beep! Time is up. It stands for Sequential Block File Manager. This is the type of data that is sent to tape drives.

Below the level of file managers are device drivers. These executable modules are real workers. They are specifically designed to work with a particular device. For every type of device there is a driver. On my Level II, Smoke Signal system, the floppy disk drive controller is their model DCB-4A. The driver is DCB4 ( good name! ). Now on my Level II, Color Computer 3 system, the driver is CC3DISK ( another good name! ). The drivers are not interchangable. They are two separate items. If by chance I got a new board for the Smoke Signal system and it was called XYZ123, the old driver would not work. And consequently the same hold true for the Color Computer System. ( I have got an disk driver for the Color Computer, Level I, that I am partial to, but it won't work with the Level II system. ) So the final word is drivers are specific. Unlike the managers which are geared to a certain class of devices, the driver is very specific to the type of device.

There are many types of devices out there. Many times they travel in groups. Do you have a one or more drives? Can your system handle more than one printer? The answer to these question is most likely, YES. And what about those extra drives or extra

printers. Are the drives one or two sided? Are they 48 tpi or 96 tpi? What about the printers? Are they running at the same baud rate? Do they both require line feeds? These questions and more are answered by the device descriptors. ( And that brings us to the this month's discussion. )

I have included an example of a device descriptor in Listing 1. I will have more to say about it a little later. But for now the listing is a good example of how to write a descriptor.

The header is similar to most modules with a few additions. The first 9 bytes consist of items like the sync bytes, module size, name offset, type, language, attributes, revision and header parity. After the parity byte comes the information that makes it a descriptor. It consists of

```
BYTE SIZE USE
$09   2   File manager offset
$0B   2   Device driver offset
$0D   1   Mode capabilities
$0E   1   Extended address
$0F   2   Port address
$11   1   Bytes in initialization table
$12   n   n bytes of initialization table
```

The initialization table follows. The one appearing in Listing 1 is for a floppy disk on my Color Computer 3, OS-9 Level II.

IT.DTP is the device type.
0...SCF
1...RBF
2...PIPE
3...SBF

IT.DRV is the drive number.
0...drive 0
1...drive 1
2...drive 2
etc.

IT.STP is the step rate. It varies with controller type. For my system the rates are

0...30mS
1...20mS
2...12mS
3....6mS

IT.TYP is the disk type

Bit 0 = 0...5 inch disk
Bit 0 = 1...8 inch disk
Bit 5 = 0...Non-color computer format
Bit 5 = 1...Color Computer format
Bit 6 = 0...Standard OS-9 format
Bit 6 = 1...Non-standard format
Bit 7 = 0...Floppy disk
Bit 7 = 1...Hard Disk

IT.DNS is the media density.
Bit 0 = 0...Single density (FM)
Bit 0 = 1...Double density (MFM)
Bit 1 = 0...Single track (48 TPI)
Bit 1 = 1...Double track (96 TPI)

IT.CYL is the number of cylinders (tracks) per side. 35, 40 and 80 per side are standard.

IT.SID is the the number of sides. It is usually 1 or 2 for floppy disks.

IT.FVY is verify disk writes.
0...verify
1...no verify

IT.SCT is sectors per track. This is usually 16 or 18 for 5 inch disk.

IT.T0S is sectors on track 0. This can vary from IT.SCT. Usual numbers are 10, 16, and 18 for 5 inch disk.

IT.ILV is the sector interleave. This is an interesting number. It is the skip factor for reading or writing a disk. After a sector is read, by the time the information is assimilated a few sectors may have passed. So this tells it how far to skip. I use a factor of $03.

IT.SAS is the minimum number of sectors allocated at one time.

### DEFAULT DRIVE

In most cases, two locations of importance are the working data directory and the execution direction. In many cases these have caused problems. If for example you wanted to hard code where you program would find some special file, you would have to settle on where the default drive would be. This was a problem with many software packages. In my C compiler, it prefers to look at /d1. What if I did not have a /d1? Too bad!

An alternate I tried is change the /d1 in the C compiler to ... ( dot-dot-dot ). This caused the compiler to look two directory levels above the current working directory. The solution was fine, except it still did not allow flexibility. Here is where the default drive comes in. It is named DD, but can be any device. It can be for drive 0 or drive 1 or a hard drive or anywhere.

In Listing 1 is a device descriptor for the Color Computer 3 using drive 0 as DD. In reality the Level II system for the Color Computer comes with a device driver DD, but the one in the Listing is different. This one has a stepping rate of 6 mS, not the 30 mS that is standard with OS-9 from Tandy.

Now it has occurred that there might be some who do not want to rewrite their device descriptors or do not have assemblers. So I created the program in Listing 2. It can accomplish the same thing, but in a different manor. This program takes an existing module and changes its internal name.

To use this one, a copy of your device descriptor must be saved. Then the program ChgName is run on it to change its name to DD. And finally the new descriptor's CRC is corrected. In terms of keyboard input, the following will work for changing D0.

    OS9: save d0
    OS9: chgname("DD") <d0 ! verify u >dd

DD will now exist and be a descriptor for drive 1. Similarily for drive 1 or a hard disk or any RBF device, a default driver can be created.

The way ChgName works is very simple. It can be broken down into steps.

1. Get the header and write it.
2. Read and write up to the name.
3. Write the new name.
4. Skip past the old name.
5. Read and write to end-of-file.

There are some problems that arise when using it to change module names . If a name is changed to something of a different length, it is possible that the new module will not work. By changing the name length, everything after it will be offset by an equivalent amount. In machine code, items that are referenced to the relative position will find that the position has changed. For example, I changed the name of DIR to D and found it crashed. On the otherhand, I changed it to my name, RON, and it worked fine. Use ChgName with discretion and caution.

Another month has come and gone. I leave you with the two programs at the end of column. I have been talking with Don Williams and hopefully in the near future, we will be offering the programs from this column on disk form. More about that later. Have a good month!

**Editor's Note:** *If you are interested in the source code for the programs that have been published, drop us a line so that I can get a fell for about how many copies to prepare.*
*Thanks*

*DMW*

## LISTING 1

```
00001          ******************************
00002          *
00003          * Name: DD
00004          * By: Ron Voigts
00005          * Date: 16-APR-88
00006          * To compile: asm dd.a o-/d0/
cmds/dd 1 #20k
00007          *
00008          ******************************
00009          *
00010          * Version 1.00  Original  RDV
00011          *
00012          ******************************
00013          *
00014          * Function:
00015          *    This is device driver used
00016          *    for a default directory.
00017          *
00018
00019                        use   /d0/
defs/defsfile
00020   0001       LEVEL     equ   1
select level one
00021              ifp1
00026              endc
00027
00028   FF40       CPort     equ   $FF40
00029
00030   00F1       TYPE      set
Devic+Object
00031   0082       REVS      set
```

ReEnt+2
```
  00032
  00033    0000 87CD0030              mod    DDEnd,DDName,TYPE,REVS,DDMgr,DDDrv
  00034
  00035          * Descriptor header
  00036    000D FF              fcb    %11111111  Device Capabilities
  00037    000E 07              fcb    $07        Extended Address
  00038    000F FF40            fdb    CPort      Disk Controler Port
  00039    0011 0F              fcb    DOptEnd-*-1
  00040
  00041          * Initialization table
  00042    0012 01       IT.DTP fcb    $01        RBF device
  00043    0013 00       IT.DRV fcb    $00        Device number 0
  00044    0014 03       IT.STP fcb    $03        6mS step rate
  00045    0015 20       IT.TYP fcb    $20        Color Computer, 5" disk
  00046    0016 01       IT.DNS fcb    $01        double densty, 48 tpi
  00047    0017 0028     IT.CYL fdb    $0028      40 tracks/side
  00048    0019 02       IT.SID fcb    $02        2 sides
  00049    001A 00       IT.VFY fcb    $00        Verify disk writes
  00050    001B 0012     IT.SCT fdb    $0012      18 sectors/track
  00051    001D 0012     IT.TOS fdb    $0012      18 sectors/track 0
  00052    001F 03       IT.ILV fcb    $03        Disk interleave of 4
  00053    0020 08       IT.SAS fcb    $08        Allocate 8 sectors opening fil
  00054    0021          DOptEnd equ   *
  00055
  00056    0021 44C4     DDName fcs    /DD/       Device name DD
  00057    0023 5242C6   DDMgr  fcs    /RBF/      Device manager
  00058    0026 43433344 DDDrv  fcs    /CC3Disk/  Device driver
  00059
  00060    002D 2A46A0          emod
  00061    0030          DDEnd  equ    *
  00062                         end
```

LISTING 2

```
PROCEDURE chgname
0000      (* *******************
0017      (*
001A      (* Name: ChgName
002A      (* By: Ron Voigts
003B      (* Date: 23-APR-88
004D      (* To compile:  Pack to CMDS directory
0073      (*
0076      (* *******************
008D      (*
0090      (* Function:
009C      (* This modudle will change an
00BA      (* executable module's name to
00D9      (* any new name.
00E9      (*
00EC      (* *******************
0103      (* Usage:
010C      (* chgname("new_name") <old_name ! verify u >new_name
0141      (* attr new_name e pe
0156      (*
0159      (* *******************
0170      (*
0173      (* External modules:
0187      (* 1. RUNB
0191      (* 2. VERIFY
019D      (* 3. ATTR
01A7      (*
01AA      (* *******************
```

```
01C1
01C2        (* New module name
01D4        PARAM new_name:STRING
01DB
01DC        (* Module Header
01EC        TYPE header=id,siz,nam:INTEGER; typ,rev,par:BYTE
020D        DIM m:header
0216
0217        (* Other variables
0229        DIM inpath,outpath:BYTE
0234        inpath:=0
023B        outpath:=1
0242
0243        DIM i:INTEGER
024A        DIM c:BYTE
0251
0252
0253        (* Get module header
0267        GET #inpath,m
0271        PUT #outpath,m
027B
027C        (* Move up to module name
0295        FOR i:=SIZE(m) TO m.nam-1
02AF          GET #inpath,c
02B9          PUT #outpath,c
02C3        NEXT i
02CE
02CF        (* Put in new name
02E1        FOR i:=1 TO LEN(new_name)
02F3          c:=ASC(MID$(new_name,i,1))
0302          IF i=LEN(new_name) THEN
0310            c:=c+$80
031C          ENDIF
031E          PUT #outpath,c
0328        NEXT i
0333
0334
0335        (* Move past old name
034A        LOOP
034C          GET #inpath,c
0356        EXITIF c>=$80 THEN ENDEXIT
0366        ENDLOOP
036A
036B        (* Process the remaider of the module
0390        WHILE NOT(EOF(#inpath)) DO
039B          GET #inpath,c
03A5          PUT #outpath,c
03AF        ENDWHILE
03B3
03B4        END


EOF
```

FOR THOSE WHO *NEED TO KNOW*

**68 MICRO JOURNAL™**

# Logically Speaking

## The Mathematical Design of Digital Control Circuits

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Mile 12 - heading for Mile 13

Well, here we are at the first anniversary of the start of this course. Happy birthday, everyone!! By way of a light introduction to Year 2, let's begin with

### COMPUTATIONAL NETWORKS

I guess you won't be TOO surprised to learn that, besides being able to control a sequential machine, relays can also be made to carry out mathematical computations. Sometimes we need to design such a circuit to demonstrate what can be done with modern design techniques, or as an attention-getter at a high-school "Look what relays can do!" type of exhibition. Examples of this type of circuit are

Given 'n', compute $n^2$ or $3n - 2$ or $n^2 - n + 1$. Or, given two numbers 'm' and 'n', compute $m * n$ or $4m - n + 1$ or $m^2 - n^2$. And so on.

In order to give you a feel for this design method, we'll develop circuits for both one and two variables. Although this sounds a LOT trickier than some of the circuits we've developed so far, it's actually not as difficult as it sounds, and we'll prove it by starting off with a simple example in a single variable.

### SINGLE-VARIABLE COMPUTATIONS

Assuming we have three pushbuttons labelled X4, X2 and X1 (corresponding to binary 4, 2 and 1), we note that by pressing various combinations (or none at all) of these buttons we can create any number from 0 through 7. The problem is to design a circuit which will square the number set up and then add 1 to it, that is, the circuit is to compute $n^2 + 1$.

| $Z \backslash^n$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---|---|---|---|---|---|---|---|---|
| L1 | 1 | | 1 | | 1 | | 1 | |
| L2 | | 1 | | 1 | | 1 | | 1 |
| L4 | | | 1 | | | | 1 | |
| L8 | | | | 1 | | 1 | | |
| L16 | | | | | 1 | 1 | | 1 |
| L32 | | | | | | | 1 | 1 |

Diagram 49

Our first step is to calculate the maximum value of the function. In this case $7^2 + 1 = 50$, which tells us that to display the result with binary-numbered lights we'll need 6 lights to indicate the binary bit-positions 32, 16, 8, 4, 2 and 1. Accordingly we construct the table of Diagram 49, with the 'n' columns headed 0 through 7, and the rows labelled L1, L2 etc up to L32. The design is carried out by calculating the value of $n^2 + 1$ for each value of 'n' given in the heading and inserting a '1' in the column below, opposite the appropriate lights. For example, in column 0, $n^2 + 1 = 0^2 + 1 = 1$, so a '1' is inserted in column 0 opposite L1. Similarly, when n = 6, $n^2 + 1 = 6^2 + 1 = 37$, and so in column 6 a '1' is inserted opposite L32, L4 and L1. Each column is systematically completed in this manner to give the final table of Diagram 49. This corresponds to the intermediate decoding-table which we'd normally construct from a flow-table.

Now, having designed VERTICALLY, we decode HORIZONTALLY to obtain our control-circuit. What could be simpler? The decodings are shown in Diagram 50. Note that our table in Diagram 49 shows quite clearly that L2 is the complement, or negation, of L1, so we merely complement L1'S decoding to obtain that for L2.

```
       4 2 1                           4 2 1
L1  0 2 4 6 | X4 X2 X1         L4   2 6 | X4 X2 X1
    x 2 4 2 | 0  0  0               x 4 | 0  1  0


       4 2 1                    4 2 1                        4 2 1
L8  3 5 | X4 X2 X1      L16  4 5 7 | X4 X2 X1      L32  6 7 | X4 X2 X1
    x   | 0  1  1            x 1 | 1  0  0              x 1 | 1  1  0
      x | 1  0  1          2 x | 1  ←  1
```

Diagram 50

L8's decoding is unusual in that not one of its variables "goes", and we are left with a 6-literal decoding, which can be reduced to five by factoring out the X1 variable.

The final decodings for the various lights are :

L1 = X1'      L2 = X1
L4 = X2.X1'      L8 = X4'.X2.X1 + X4.X2'.X1 = X1(X4'.X2 + X4.X2')
L16 = X4.X2' + X4.X1 = X4(X2' + X1)      L32 = X4.X2

Of course, if we wished instead to have the result displayed decimally, we would have coded vertically for L1, L2, L3 ... L8, L9, L10, L20, L30, L40 and L50.

MULTI-VARIABLE COMPUTATIONS

Now for a SLIGHTLY more difficult example! Let's assume we have TWO sets of pushbuttons, enabling us to set up values for the number 'm' and the number 'n', and our circuit has to multiply these two numbers together, double the result and then add 1, that is, it must compute 2mn + 1. We'll work this one out with the added restriction that both 'm' and 'n' are limited to the values 0 through 3.
It's obvious that both numbers will require 2 buttons, with the binary values 1 and 2, in order to set up the figures 0 to 3, giving a total of 4 buttons. Four buttons means a total of 4 binary bits, which, in turn, are capable of representing the combined binary range 0 through 15. This tells us that our intermediate-table will require 16 columns, headed 0 to 15, with an additional 2 rows of data above this header - one for 'n', and one for 'm' above that.
The highest number we can compute is when both 'm' and 'n' are at their maximum, that is 3, giving us 2 * 3 * 3 + 1 = 19, for which we're going to need output lights ranging from L1 through L16, or 5 output rows.

| m | 0 | | | | 1 | | | | 2 | | | | 3 | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| n | 0 | 1 | 2 | 3 | 0 | 1 | 2 | 3 | 0 | 1 | 2 | 3 | 0 | 1 | 2 | 3 |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| L1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| L2 |  |  |  |  |  | 1 |  | 1 |  |  |  |  |  | 1 |  | 1 |
| L4 |  |  |  |  |  |  | 1 | 1 |  | 1 |  | 1 |  | 1 | 1 |  |
| L8 |  |  |  |  |  |  |  |  |  |  | 1 | 1 |  |  | 1 |  |
| L16 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 1 |

Diagram 51

Note that we set up the heading for 'n' by repeating the sequence 0 - 3 for each of the values 0 - 3 for 'm'. We have thus paired off all possible combinations of 'm' with all possible combinations of 'n'.
As before, we design vertically, commencing with column 0, where both 'm' and 'n' are equal to 0, so that 2mn + 1 = 1. We record this by by placing a '1' opposite L1 in Column 0, and similarly in the next 3 columns where m = 0. In Column 10, where both 'm' and 'n' equal 2, 2mn + 1 equals 9, which is recorded by placing a '1' opposite L1 and L8 in this column. Once the whole table is completed, we decode horizontally as before, noting, by the way, that the row for L1 is completely filled with 1s. This indicates that L1 is permanently ON from the moment of switching ON our circuit - it's permanently

wired across the power supply.

I'll leave the actual decoding to you as an additional exercise, as the method is quite straightforward (no phis in this table) and the extra practice will not come amiss. Don't skip doing this, though you may use your decoder-cards if you like! I'll just give you one piece of info, and that is that the decoding-table's variable header will look like

    8 4 2 1
    m2 m1 n1 n2

Notice that 'm' comes first, because it occupies the higher level of the intermediate-table's header, and is then followed by 'n' - - with two binary bits apiece.

## CIRCUIT IMPLEMENTATION WITH STANDARD AND-OR-NOT LOGIC UNITS

Normally I wouldn't be dealing with solid-state for a little while to come, as I used to give these lessons on a weekly basis, but as we're already into our 13th month I think I'll take a little time out to show you how to implement AND-OR-NOT logic units in our circuit design. I'll leave NOR and NAND implementation till later, as the decoding technique for these is a LITTLE more complex.

The whole design technique for standard, or AND-OR-NOT, logic is exactly the same as for relay implementation. The one big difference is that we won't need multi-pole pushbuttons, or even buttons with complementary contacts. A simple light-weight button with only a single NO-contact will suffice. We'll use our single-variable example of $n^2 + 1$ to show how the various logic units are connected.
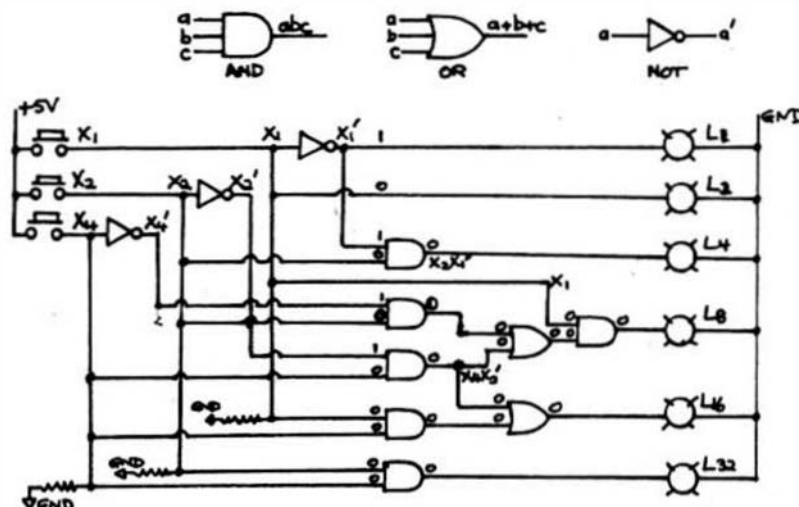


Diagram 52

Along the top of Diagram 52 the logic symbols for the three devices are shown. For the enlightenment of newcomers to this field, I'll explain that an AND unit will output a binary '1' if all its inputs are equal to '1' (that is, input-A AND input-B AND input-C, etc, are all equal to '1'), whereas an OR unit will output a logic '1' if either input-A OR input-B OR input-C, etc, is equal to '1' (OR any combination thereof), otherwise the output is equivalent to logic-0. A NOT unit merely outputs the complement of its input, outputting a '1' if the input is '0', and vice-versa.

It's fairly standard practice to keep the inputs at the left, and to work towards the outputs at the right. We may also use any intermediate function of one of the outputs as part of another output's circuit. For example, the X4.X2 of L8 can be used for L16, though this wouldn't be possible with relays. In addition, as we're now working with low-voltage circuitry, we could use LEDs for our output lights.

And so, back to relays for a while.

## DESIGNING INVERSE NETWORKS

Suppose we have a circuit controlling L1, as shown in Diagram 53a, and we desire the inverse, or complementary, network to control L2. That is to say, we need a network which will guarantee that whenever L1 is ON, L2 will be OFF, and whenever L1 is OFF then L2 will be ON.

There are several ways of designing such a circuit. We could, for instance, merely negate the Boolean expression for the original network, and draw the circuit-diagram for the resultant expression. The present network is expressed as

$$L1 = (a + c')(b' + de') + f$$

and its complement is therefore

$$L2 = [(a + c')(b' + de') + f]' = (a'c + b(d' + e))f$$

Another method would be to enter the expression on a K-map in the form of 1s, and then read out the 0s as if they were 1s. But as we haven't yet learned advanced K-mapping, we can't for the moment comfortably handle more than 4-variable maps. There is another complication, too, in that, to date, we've been dealing only with ladder-type relay networks. That is, we've constructed only series-parallel networks, and the above techniques would be almost useless, perhaps even worthless, in the case of "bridge" or other non-series-parallel networks. However, by using "graphical complementation" all networks, with the exception of non-planar networks, CAN be very easily complemented.
A "non-planar" network, by the way, is one which CANNOT be drawn on a sheet of paper without at least one connecting-line crossing another. Several networks you may come across may LOOK like non-planar ones, but very often the line-crossings can be eliminated by re-arranging the drawing. Only if it's ABSOLUTELY IMPOSSIBLE to do so is the network truly non-planar. We'll be dealing with this type of network in the not-too-distant future in a lot more detail!
Anyway, returning to 53a, the first step in obtaining the inverse network is to place a dot, or "node", inside each closed loop, or "mesh", of the network, as shown in 53b. In addition, an extra node is placed above the network and one below. Now, if it's possible to get from one node to another by passing through ONE contact only, then a dotted line should be drawn connecting these two nodes. All such possibilities are indicated in 53c, and L2 has been added to the lower end. For the sake of clarity, the contacts have not been labelled in this diagram.
If we now re-draw this diagram, as in 53d, AND COMPLEMENT ALL THE CONTACTS, we've drawn the complementary network called for in our specifications. Of course, the contact network of 53d would actually be swung round anti-clock-wise to form the normal right-to-left ladder configuration.
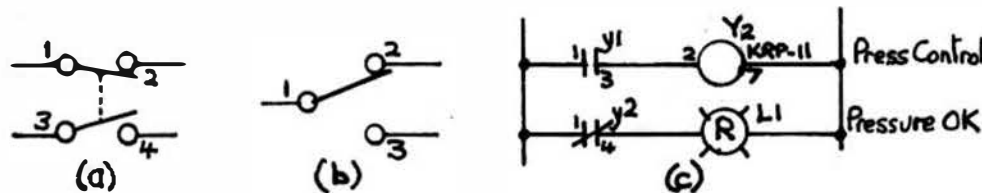
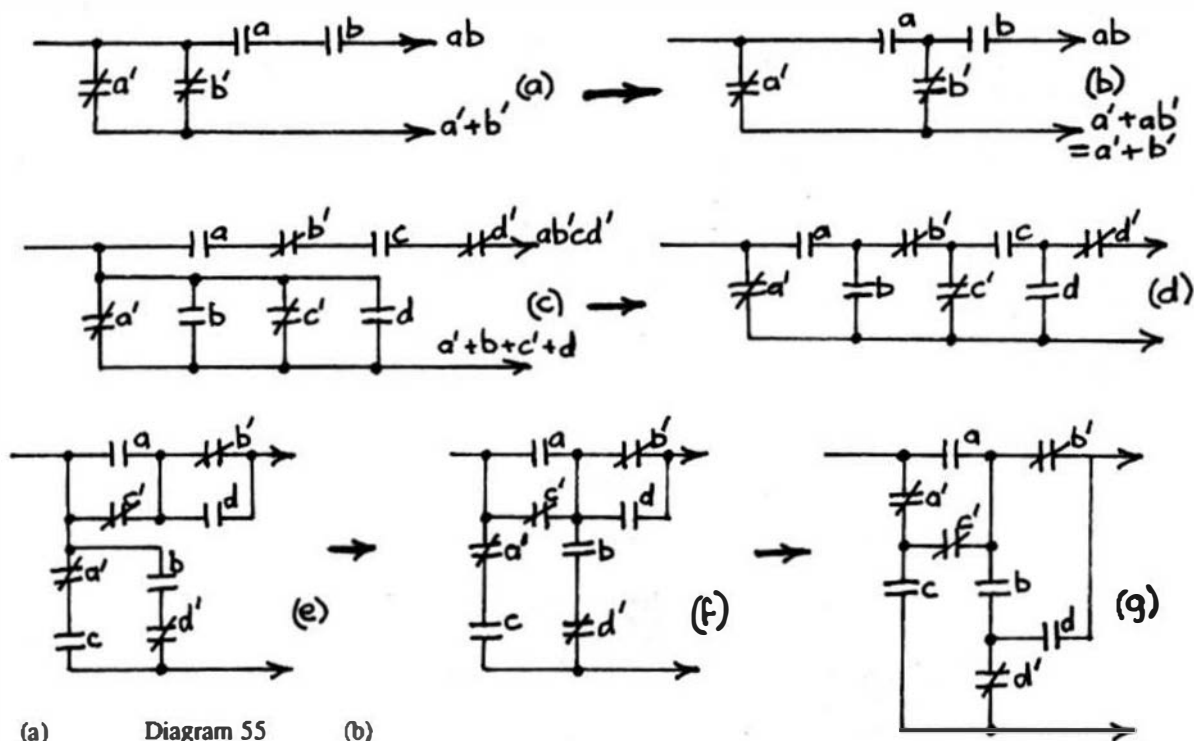## TRANSFER-CONTACT NETWORKS



Diagram 54

Relay contacts may come in different configurations, Diagram 54 showing two of the most common. In 54a we see separate NO (normally-open) and NC (normally-closed) contacts, and, of course, there may be several such pairs on a single relay. Diagram 54b shows a combined NO and NC configuration, called a "transfer-contact" because the switch transfers the current-flow from the upper to the lower contact.
Observe the numbers attached to the contacts. It's normal procedure with plug-in type relays to incorporate the relay's pin-numbers into the circuit-diagram, with the left-most number usually above the horizontal line and the right-hand number below, wherever possible. For larger, hard-wired relays we'll designate the various contacts as y1a, y1b, and so on, to ease the burden of trouble-shooting later. Note also in 54c, that the relay-coil not only includes its pin-numbers, but has the relay-type or catalog-number added, together with a brief description of its function. Lights will have an internal letter denoting their colour, "R" for red, "G" for green, etc, plus a description of their function.
Transfer-contacts occur so frequently that it's well worthwhile looking at a few techniques specially directed towards minimising the number of "transfers", as they're usually called, where this type of relay is utilised. Let's see what can be done, assuming that we're using only transfer relays.

# 1. COMBINING A NETWORK AND ITS COMPLEMENT



(a)　　　　Diagram 55　　　　(b)

If we have a simple network with two outputs, one of which is the exact complement of the other, as in 55a, the circuit as shown would need 3 transfers - one on relay-A (because the a and the a' contacts have a common point) and two on relay-B, using the NO-contact of one transfer and the NC-contact of the second. However, if we break apart the section of circuit composed of contacts in parallel, and transfer the "loose end" to the appropriate part of the series path, the function will remain unchanged. In 55a we've lifted off the wire joining one end of the b'-contact to the a-contact, and transferred it to the junction of the a-contact and the b-contact in the series path, thus giving both contacts on relay-B a common point too. This move reduces the total number of transfers from 3 to 2, without changing the function of either path.

Two more examples are given in 55c and 55d, and 55e, f and g. In the last example, note how the process of breaking-up and re-arranging the parallel paths is repeated to produce a network which is more difficult to understand, but which, nevertheless, has reduced the total number of transfers required from seven to four.

Don't forget, this method is used WHERE ONE NETWORK IS THE COMPLEMENT OF THE OTHER.

## 2. COMBINING LIKE CONTACTS IN SEPARATE PARALLEL PATHS

Sometimes it turns out that more CONTACTS will be required in a network in order to reduce the number of TRANSFERS, as we'll see next.



Diagram 56

Diagram 56 shows how the b-contact (which really symbolises ANY network common to the two outputs) need be used only once, by inserting the complement of the other contact into the network. Each of the original parallel paths is thus transformed from the form a + b into its equivalent form a + a'b (remember Rule 6 of The Laws of Boolean Algebra?).

The second example demonstrates the use of both this technique and the previous one to obtain a minimum transfer realisation of the initial network.

In both these cases we end up with MORE contacts, but a lesser number of transfers, which does represent an actual saving, as the complementary contacts were already available AND UNUSED on the relays concerned.

Another technique is to use a DC source of power (instead of AC) as we've assumed till now). Then, by utilising DC relays, a couple of diodes will do the trick instead of figuring out a re-arranged contact network. The diodes, of course, serve to isolate the outputs so that the remainder of one's network will not activate the other output.

If, as sometimes happens, the parallel paths contain a contact and its complement in addition to the common like contact (or block of contacts), the common like contacts can be shifted so that a bridge is formed between the two outputs, with no change in either of the functions.



Diagram 57

Diagram 57 illustrates the method for two different examples, and shows how a common block, no matter how large, need be constructed only once PROVIDED ONE OF THE OUTPUT PATHS CONTAINS THE COMPLEMENT OF ONE BRANCH OF THE OTHER OUTPUT PATH. Again you should note that although only a single complementary contact is shown in each example, it's symbolic of ANY path, no matter how complicated.

Diagram 58 gives another example of how it's possible to use more than one of these techniques to re-arrange a network.



Diagram 58

## RELAY TREES

A relay tree, or transfer-tree, has a single power input and any number of distinctly separate outputs, such that each output path passes through one contact on each relay, and only one output is energised at any one time. Such a tree is shown in Diagram 59, where each of the eight separate outputs is energised through one contact on each of the three relays A, B and C. The fact that only one output is energised at any one time is easily verified by converting the Boolean expression for each path into binary, and thence into decimal if need be. In our diagram, the outputs represent the binary numbers 000 to 111 inclusively, or the decimal numbers 0 to 7.

Diagram 59

Notice that while only one transfer-contact is used on relay A, two are necessary on relay-B and four on relay-C. This, of course, is the binary sequence 1, 2, 4 ..... and means that if we had a 6-relay tree the sixth relay would require 32 transfer contacts. If the number of relays is represented by 'n', the number of transfer-contacts necessary on the nth relay would be $2^{(n-1)}$. That is, in the case of six relays it would be $2^5$, or 32.

A full relay-tree has an output for each and every possible relay combination, and is composed of a total of $2^n - 1$ transfers (NOT the same as our previous formula). In the case of a 6-relay tree this works out to $2^6 - 1 = 64 - 1 = 63$ transfer-contacts altogether!

Returning to Diagram 59, the transfer-contact distribution is

A B C
1 2 4 = 7 transfers

Although it's not possible to reduce the TOTAL number of transfers in a FULL relay-tree, the distribution of the contact-load CAN be equalised by re-arranging the tree, apart from the first relay connected to the power-input - which has only one transfer.



Diagram 60

Diagram 60 shows how the remaining 6 transfers are divided equally between relays B and C to use only 3 transfers apiece. The distribution is now

A B C
1 3 3 = 7 transfers

If the relay-tree is reversed from left-to-right so that we have any number of distinct inputs connected to only one output, the network is known as a reverse tree.

## PARTIAL TREES

Now, while the total number of transfers cannot be reduced in a FULL relay-tree, it's possible to do so in the case of a PARTIAL tree by re-arranging the tree. A partial tree is one in which some of the "branches" are missing. Moreover, there's a technique, which uses a K-map in a completely different way from heretofore, whereby we can determine the BEST method of re-arranging the transfers.

Let's look at an example to see how it's done. Suppose we're called on to minimise the following output combinations in relay-tree form, keeping in mind that each term represents a SEPARATE output

$a'b'c'd'$   $a'b'c'd$   $a'bc'd'$   $a'bc'd$   $ab'cd'$   $ab'cd$   $abc'd'$   $abcd'$

Normally, a K-map is used to record the function for only one output, but in the case of relay-trees each 1-entry corresponds to a single output of the tree, AND MUST THEREFORE BE READ AS A SINGLE ENTRY. There must be no forming of loops, as this would be equivalent to bundling a set of independent output functions as though they were one.



Diagram 61

Our first step, Diagram 61a, is to enter 1s on the K-map to correspond with each of the required outputs. The design process consists of cutting the K-map into two EQUAL parts, each of which is called a "sub-map", and then to cut these sub-maps into two EQUAL parts also, to form new sub-maps, in such a way that the UNUSED squares of the K-map form AS FEW and AS LARGE groups as possible.

## DISSECTING THE K-MAP

For example, we could cut map 61a into two with a vertical slice down the middle, and then by cutting the left-hand submap horizontally into two we'd isolate the unused squares marked with 'x'. Next we'd cut our right-hand sub-map into two submaps by a vertical slice, and a further horizontal slice in the extreme right-hand sub-map would isolate the block marked 'z'. Finally, and this is a little trickier, we recall that that the top and bottom of a K-map are really adjacent, and so, by imagining our original column 11 as a continuous band, we can make two quick snips horizontally to isolate the block 'y' from the two 1s, which would still be joined together. We'd thus end up with three blocks of unused contacts, so this will be our approach. Of course, we don't ACTUALLY cut the map, but mark it, maybe in red, where our cuts are going to made!

Suppose we were to make our first cut horizontally instead of vertically as we did above! You can see that it would go right through the middle of the two ys, and the best that we could hope for after that would be FOUR blocks of unused squares, consisting of our original 'x' and 'z' and two small blocks of 'y'.

Following our first plan then, we'll make a vertical cut down the centre of the K-map, shown by the heavier line in 61b. If we look carefully at the column headings we see that what we've done is to separate all squares which contain a' as a component from those which contain a. In other words, we've created a a'-submap and an a-submap, and we record this fact by making an appropriate entry in each of the 1-squares, as shown in the diagram.

Proceeding to 61c, we now make a horizontal cut in the a'-submap, and record in its 1-squares the fact that we've thereby separated the cs from the c's. Still with 61c, we'll cut the a-submap vertically into two, forming an ab-submap and an ab'-submap. This, too, is recorded in the corresponding 1-squares.

Moving now to 61d, we make two snips in the "band" forming the ab-submap (remember the 1s are adjacent, and therefore joined), thus isolating an unused block and forming an abd'-submap of 1s. We'll also chop the ab'-submap into two, to give a final isolated block of unused squares and an ab'c-submap of 1s.

Having reached the stage where ALL the unused blocks have been isolated, the last step of the procedure is to concentrate on the submaps containing 1-entries, and to chop them up by the same technique as before until each 1-square is a complete submap in its own right. Diagrams 61e and 61f show these final stages, with the literals being recorded in the various squares as the splitting process is carried out.



Diagram 62

The relay-tree is constructed from the K-map of 61f, as shown in Diagram 62. The method is to take the 1-squares of the K-map one at a time, commencing in the top-left corner, that is, square 0000. A path is then drawn from power-input to output containing contacts in the EXACT order shown in this square, ie, a' c' d' and b'. The square to the right (0100) has the entry a' c' d' b, but as the first three contacts have already been drawn we merely branch off with a b-contact from the junction of d' and b'. Similarly, when we come to square 0001, we find that we already have the a'c' part, so we branch off into the db' section. And so on for the remaining squares, which gives us a transfer distribution of

    A B C D
    1 3 3 3 = 10 transfers

And here is where I think we should take a little break, so as not to monopolise TOO much of this magazine, and as you've had things fairly easy of late I'm going to give you another test to sharpen your skills. Here then is

TEST TEN

1. Design the following computational networks, assuming 'n' to range over the values 0 - 7

    (a) $n^2$    (b) $3n + 1$    (c) $2 * n^2 - n + 1$

2. Design the following computational networks, assuming 'm' to range over the values 0 - 3, and 'n' over the values 0 - 5

    (a) $2m + n$    (b) $3m + 2n$    (c) $3mn + 2$

3. Complement the following networks using the graphical complementation technique

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* Now supplied as a two disk set:

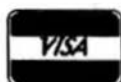*Disk #1: JUST2.CMD object file,*

*JUST2.TXT PL9 source:FLEX, SK*DOS - CC*

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source compiles to a standard syntax JUST.CMD object file. Using JUST syntax (.p ,u ,y etc.) With all JUST functions plus several additional printer formatting functions. Reference the JUSTSC C source. For those wanting an excellent BUDGET PRICED word processor, with features none of the others have. This is it!

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   3.  DPOKE -- Double poke
   4.  FPOS -- Current file position
   5.  FSIZE -- File size
   6.  FTRIM -- removes leading spaces from a string
   7.  GETPR -- returns the current process ID
   8.  GETOPT -- gets 32 byte option section
   9.  GETUSR -- gets the user ID
   10. GTIME -- gets the time
   11. INSERT -- insert a string into another
   12. LOWER -- converts a string into lowercase
   13. READY -- Checks for available input
   14. SETPRIOR -- changes a process priority
   15. SETUSR -- changes the user ID
   16. SETOPT -- set 32 byte option packet
   17. STIME -- sets the time
   18. SPACE -- adds spaces to a string
   19. SWAP -- swaps any two variables
   20. SYSCALL -- system call
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   **Limited Special - $19.95**

## SOFTOOLS

The following programs are included in object form for immediate application. PL/9 source code available for customization.

**READ-ME** Complete instructions for initial set-up and operation. Can even be printed out with the included text processor.

**CONFIG** one time system configuration.

**CHANGE** changes words, characters, etc. globally to any text type file.

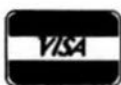**CLEANTXT** converts text files to standard FLEX, SK*DOS files.

**COMMON** compare two text files and reports differences.

**COMPARE** another check file that reports mis-matched lines.

**CONCAT** similar to FLEX, SK*DOS append but can also list files to screen.

**DOCUMENT** for PL/9 source files. Very useful in examining parameter passing aspects of procedures.

ECHO echos to either screen or file.

FIND an improved *find* command with "*pattern*" matching and wildcards. Very useful.

HEX dumps files in both hex and ASCII.

INCLUDE a file copy program that will accept "includes" of other disk files.

KWIC allows rotating each word, on each line to the beginning. Very useful in a sort program, etc.

LISTDIR a directory listing program. Not super, but better than CAT.

MEMSORT a high-speed text file sorter. Up to 10 fields may be sorted. Very fast. Very useful.

MULTICOL width of page, number of columns may be specified. A MUST!

PAGE similar to LIST but allows for a page header, page width and depth. Adjust for CRT screen or printer as set up by CONFIG. A very smart print driver. Allows printer control commands.

REMOVE a fast file deleter. Careful, no prompts issued. Zap, and its gone!

SCREEN a screen listing utility. Word wraps text to fit screen. Screen depth may be altered at run time.

SORT a super version of MEMSORT. Ascending/descending order, up to 10 keys, case over-ride, sort on n° word and sort on characters if file is small enough, sorts in RAM. If large file, sort is constrained to size of your largest disk capacity.

TPROC a small but nice text formatter. This is a complete formatter and has functions not found in other formatters.

TRANSLIT sorts a file by x keyfields. Checks for duplications. Up to 10 key files may be used.

UNROTATE used with KWIC this program reads an input file and unfolds it a line at a time. If the file has been sorted each word will be presented in sequence.

WC a word count utility. Can count words, characters or lines.

NOTE: this set of utilities consists of 6 5-1/4" disks or 2 8" disks, w/ source (PL9). 3 5-1/4" disks or 1 8" disk w/o source. Complete set SPECIAL INTRO PRICE:
5-1/4" w/source FLEX - SK*DOS - $129.95
w/o source - $79.95
8" w/source - $79.95 - w/o source $49.95

FULL SCREEN FORMS DISPLAY from Computer Systems Consultants - TSC Extended BASIC program supports any Serial Terminal with Cursor Control or Memory-Mapped Video Displays; substantially extends the capabilities of the Program Designer by providing a table-driven method of describing and using Full Screen Displays.
*F, S and CCF, U - $25.00, w/ Source - $50.00*

SOLVE from S.E. Media - OS-9 Levels I and II only. A Symbolic Object/Logic Verification & Examine debugger. Including inline debugging, disassemble and assemble. SOLVE IS THE MOST COMPLETE DEBUGGER we have seen for the 6809 OS-9 series! SOLVE does it all! With a rich selection of monitor, assembler, disassembler, environmental, execution and other miscellaneous commands, SOLVE is the MOST POWERFUL tool-kit item you can own! Yet, SOLVE is simple to use! With complete documentation, a snap! Everyone who has ordered this package has raved! See review - 68 Micro Journal - December 1985. No 'blind' debugging here, full screen displays, rich and complete in information presented. Since review in 68 Micro Journal, this is our fastest mover!
*Levels I & II only - OS-9 $69.95*

## DISK UTILITIES

OS-9 VDisk from S.E. Media -- For Level I only. Use the Extended Memory capability of your SWTPC or Gimix CPU card (or similar format DAT) for FAST Program Compiles, CMD execution, high speed inter-process communications (without pipe buffers), etc. - SAVE that System Memory. Virtual Disk size is variable in 4K increments up to 960K. Some Assembly Required.
*Level I OS-9 obj. $79.95; w/ Source $149.95*

O-F from S.E. Media -- Written in BASIC09 (with Source), includes: REFORMAT, a BASIC09 Program that reformats a chosen amount of an OS-9 disk to FLEX, SK*DOS Format so it can be used normally by FLEX, SK*DOS; and FLEX, a BASIC09 Program that does the actual read or write function to the special O-F Transfer Disk; user-friendly menu driven. Read the FLEX, SK*DOS Directory, Delete FLEX, SK*DOS Files, Copy both directions, etc. FLEX, SK*DOS users use the special disk just like any other FLEX, SK*DOS disk
*O - 6809/68000 $79.95*

LSORT from S.E. Media - A SORT/MERGE package for OS-9 (Level I & II only). Sorts records with fixed lengths or variable lengths. Allows for either ascending or descending sort. Sorting can be done in either ASCII sequence or alternate collating sequence. Right, left or no justification of data fields available. LSORT includes a full set of comments and errors messages.
*OS-9 $85.00*

HIER from S.E. Media - *HIER is a modern hierarchal storage system for users under FLEX, SK*DOS.* It answers the needs of those who have hard disk capabilities on their systems, or many files on one disk - any size. Using HIER a regular (any) FLEX, SK*DOS disk (8 - 5 - hard disk) can have sub directories. By this method the problems of assigning unique names to files is less burdensome. Different files with the exact same name may be on the same disk, as long as they are in different directories. For the winchester user this becomes a must. Sub-directories are the modern day solution that all current large systems use. **Each directory looks to FLEX, SK*DOS like a regular file,** except they **have the extension '.DIR'.** A full set of directory handling programs are included, making the operation of HIER simple and straightforward. A special install package is included to install HIER to your particular version of FLEX, SK*DOS. Some assembly required. Install indicates each byte or reference change needed. Typically - 6 byte changes in source (furnished) and one assembly of HIER is all that is required. No programming required!
*FLEX - SK*DOS $79.95*

COPYMULT from S.E. Media -- Copy LARGE Disks to several smaller disks. FLEX, SK*DOS utilities allow the backup of ANY size disk to any SMALLER size diskettes (Hard Disk to floppies, 8" to 5", etc.) by simply inserting diskettes as requested by COPYMULT. No fooling with directory deletions, etc. COPYMULT.CMD understands normal "copy" syntax and keeps up with files copied by maintaining directories for both host and receiving disk system. Also includes BACKUP.CMD to download any size "random" type file; RESTORE.CMD to restructure copied "random" files for copying, or recopying back to the host system; and FREELINK.CMD as a "bonus" utility that "relinks" the free chain of floppy or hard disk, eliminating fragmentation.
*Completely documented Assembly Language Source files included. ALL 4 Programs (FLEX, SK*DOS, 8" or 5") $99.50*

COPYCAT from Lucidata -- *Pascal NOT required.* Allows reading TSC Mini-FLEX, SK*DOS, SSB DOS68, and Digital Research CP/M Disks while operating under SK*DOS, FLEX1.0, FLEX 2.0, or FLEX 9.0 with 6800 or 6809 Systems. COPYCAT will not perform miracles, but, between the program and the manual, you stand a good chance of accomplishing a transfer. Also includes some Utilities to help out. Programs supplied in Modular Source Code (Assembly Language) to help solve unusual problems.

    F, S and CCF 5" - $50.00     F, S 8" - $65.00

VIRTUAL TERMINAL from S.E. Media - Allows one terminal to do the work of several. The user may start as many as eight tasks on one terminal, under *VIRTUAL TERMINAL* and switch back and forth between tasks at will. No need to exit each one; just jump back and forth. Complete with configuration program. The best way to keep up with those background programs.

    6809 O & CCO - obj. only - $49.95

FLEX, SK*DOS DISK UTILITIES from Computer Systems Consultants -- Eight (8) different Assembly Language (w/ Source Code) FLEX, SK*DOS Utilities for every FLEX, SK*DOS Users Toolbox: Copy a File with CRC Errors; Test Disk for errors; Compare two Disks; a fast Disk Backup Program; Edit Disk Sectors; Linearize Free-Chain on the Disk; print Disk Identification; and Sort and Replace the Disk Directory (in sorted order). -- PLUS -- Ten XBASIC Programs including: A BASIC Resequencer with EXTRAs over "RENUM" like check for missing label definitions, processes Disk to Disk instead of in Memory, etc. Other programs Compare, Merge, or Generate Updates between two BASIC Programs, check BASIC Sequence Numbers, compare two unsequenced files, and 5 Programs for establishing a Master Directory of several Disks, and sorting, selecting, updating, and printing paginated listings of these files. A BASIC Cross-Reference Program, written in Assembly Language, which provides an X-Ref Listing of the Variables and Reserved Words in TSC BASIC, XBASIC, and PRECOMPILER BASIC Programs.

ALL *Utilities include Source (either BASIC or A.L. Source Code).*
    F, S and CCF - $50.00
    BASIC *Utilities ONLY for UniFLEX* -- $30.00

MS-DOS-FLEX Transfer Utilities to OS-9 For 68XXX and CoCo* OS-9 Systems Now READ - WRITE - DIR - DUMP - EXPLORE FLEX & MS-DOS Disk. These Utilities come with a rich set of options allowing the transfer of text type files from/to FLEX & MS-DOS disks. *CoCo systems require the D.P. Johnson SDISK utilities and OS-9 and two drives of which one must be a "host" floppy.
    *CoCo Version: $69.95     68XXX Version $99.95

## MISCELLANEOUS

TABULA RASA SPREADSHEET from Computer Systems Consultants -- TABULA RASA is similar to DESKTOP/PLAN; provides use of tabular computation schemes used for analysis of business, sales, and economic conditions. Menu-driven; extensive report-generation capabilities. Requires TSC's Extended BASIC.
    F, S and CCF, U - $50.00, w/ Source - $100.00

DYNACALC -- Electronic Spread Sheet for the 6809 and 68000.
    F, S, OS-9 and SPECIAL CCF - $200.00, U - $395.00
    OS-9 68K - $595.00

FULL SCREEN INVENTORY/MRP from Computer Systems Consultants Use the Full Screen Inventory System/Materials Requirement Planning for maintaining inventories. Keeps item field file in alphabetical order for easier inquiry. Locate and/or print records matching partial or complete item, description, vendor, or attributes; find backorder or below stock levels. Print-outs in item or vendor order. MRP capability for the maintenance and analysis of Hierarchical assemblies of items in the inventory file. Requires TSC's Extended BASIC.
    F, S and CCF, U - $50.00, w/ Source - $100.00

FULL SCREEN MAILING LIST from Computer Systems Consultants -- The Full Screen Mailing List System provides a means of maintaining simple mailing lists. Locate all records matching on partial or complete name, city, state, zip, or attributes for Listings or Labels, etc. Requires TSC's Extended BASIC.
    F, S and CCF, U - $50.00, w/ Source - $100.00

DIET-TRAC Forecaster from S.E. Media -- An XBASIC program that plans a diet in terms of either calories and percentage of carbohydrates, proteins and fats (C P G%) or grams of Carbohydrate, Protein and Fat food exchanges of each of the six basic food groups (vegetable, bread, meat, skim milk, fruit and fat) for a specific individual. Sex, Age, Height, Present Weight, Frame Size, Activity Level and Basal Metabolic Rate for normal individual are taken into account. Ideal weight and sustaining calories for any weight of the above individual are calculated. Provides number of days and daily calendar after weight goal and calorie plan is determined.
    F, S - $59.95, U - $89.95

## GAMES

RAPIER - 6809 Chess Program from S.E. Media -- Requires FLEX, SK*DOS and Displays on Any Type Terminal. Features: Four levels of play, Swap side. Point scoring system. Two display boards. Change skill level. Solve Checkmate problems in 1-2-3-4 moves. Make move and swap sides. Play white or black. *This is one of the strongest CHESS programs running on any microcomputer, estimated USCF Rating 1600+ (better than most 'club' players at higher levels)*
    F, S and CCF - $79.95

### NEW

MS-DOS/FLEX *Transfer Utilities* For 68XXX and CoCo* OS-9 Systems. Now Read, Write, DIR, Dump and Explore FLEX & MS-DOS Disks. Supplied with a rich set of options to explore and transfer text type files from/to FLEX and MS-DOS disks. *CoCo OS-9 requires SDISK utilities & two floppy drives.
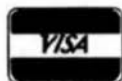CCO $69.95     68XXX OS-9 $99.95

## *Macintosh Software at Discounted Prices*
### *"Call for prices, it"ll be worth the savings."*

NOTE: Changes

1   Price increase for SCULPTOR, see advertising front of this catalog and other ad in this issue. Special price for 68 Micro Journal readers.

2.   Lower price for BASIC09 TOOLS, see Utilities section.

3.   New MS-DOS & FLEX to OS-9 Utilities, see above.

(a)    (b)    (c)    (d)

Observe that networks (c) and (d) are not our normal ladder type. For instance, in (c), is the first loop "a + ch" OR maybe "c + ah"? We'll be designing networks like this, which cannot be described in ordinary Boolean algebra, a bit later on.

4. Design the following networks, using transfer contacts as economically as possible.

(a) $L1 = abc + a'b'c'$
(b) $L1 = a(b + c') + de'$    $L2 = (a' + b'c)(d' + e)$
(c) $L1 = ab + c$        $L2 = ad + e$
(d) $L1 = a + bc$        $L2 = a + d + e$
(e) $L1 = a(b + c) + d + e'$    $L2 = a(b + c) + e + f'$

5. Equalise the contact-load on a full relay-tree of four variables a, b, c and d, and draw the final tree.

6. Minimise the number of transfer-contacts in the following partial trees

(a) a'b'c'd'  a'b'cd'  a'b'cd  a'bc'd  a'bcd  ab'c'd'  abcd'  abcd
(b) a'bc'd'  a'bcd'  ab'c'd  ab'cd  abc'd'  abc'd  abcd'  abcd
(c) abc'd'  abc'd  ab'c'd'  a'bcd  a'b'cd  a'b'c'd  ab'cd'  abcd

It sometimes happens that it's impossible to obtain the APPARENT unused groupings by the sub-dividing process, and a group just HAS to be broken up. Exercise 6c is a good example of this condition! Don't forget that each map, or sub-map, must be divided into two EQUAL parts at each cutting operation. You CANNOT cut a 4-square submap into a 1-square and a 3-square, for instance. Anyway, have fun with this lot!!

... End of Mile 12, sitting at marker 'Mile 13'

+++

*FOR THOSE WHO NEED TO KNOW*    68 MICRO JOURNAL™

# A Review of SideKick 2.0

## Desktop Organizing Software From Borland, International

By: James E. Law
1806 Rock Bluff Rd.
Hixson, TN, 37343

A number of software packages advertise that they can get us organized. Well, I for one, need to get organized so I was interested in trying the new version of Borland's SideKick.

This package has been around a long time, both for the Macintosh and IBM compatibles. SideKick contains a collection of desktop organization programs including telecon aids (e.g., dialing, phone books, and logs), calendars, a clock, a calculator, an enhanced notepad, a print spooler, and a simple specialized data base. In addition, the latest version (2.0) adds outlining and a spreadsheet.

SideKick applications, which comes on two 800k disks, can be used on any Macintosh. Several of the programs are provided in both application and desk accessory formats so that they can be used with current versions of MultiFinder.

With preliminaries aside, let's look at each of SideKick's major subprograms.

### First, Make an Outline

Remember in high school English classes when you were admonished to prepare an outline before writing your essay? If you had had your Macintosh and SideKick back then, the outline would have been no sweat. SideKick's outliner, called Outlook, is capable and relatively full featured. It may be all the outlines most users ever need.

Upon selecting Outlook from the Apple menu, you are given the option of opening an existing outline, starting a new outline, or resuming the last outline opened. Multiple outlines may be open at the same time.

Each entry in an outline begins with a small circle (i.e., a "bullet"). The bullet for a major heading is hollow if there are no lower tier entries under that heading. Otherwise, it is solid black. The bullet is followed by an icon that looks like a miniature page. This icon changes to indicate when there is text or graphics documents filed under the related heading. You may choose "Hide Icons" from the Outlook menu to get rid of the page icons.

To begin an outline you merely place the cursor to the right of the icon and begin typing. When you press return, a new bullet and page icon are entered and the cursor is automatically placed for your new entry. Any entry can be selected by clicking on its bullet. Selected entries can be moved to the right or left (to change their relative rank) by using the cursor keys or by dragging the entry with the mouse. Selected entries can also be deleted, copied, or cut.

Each entry on your outline may be associated with a text or graphics documents. By clicking on the page icon, this document can be opened, viewed, copied, cut, or (in the case of text) modified. For example, suppose that your outline was a list of actions needed to get ready for a presentation.

It might look like this:

- Preparation for sales presentation meeting
  - Prepare agenda
  - Prepare advertising material for review
    - Script for radio spot
    - Billboard layout
- Prepare overhead transparencies
  - Title transparency
  - Sales-by-district transparency
- Obtain meeting room

When you click the page icon for the entry entitled "script for radio spot," a text-only file would be opened with your script. Similarly, clicking the icon for "sales-by-district" might display a bar draft of sales statistics. You can see how an outline can be used as a powerful special-purpose data base including both text and graphics.

A specific document under an Outlook entry is either treated as text or graphics but not both. Also, text files are text only and do not include special formating information (i.e., multiple fonts and styles).

The Page Layout option from the Outlook menu allows you considerable flexibility in determining how your outline will look when it prints. For example, you can specify margins, line spacing, how entries will be numbered (e.g., 1, 1.1, 1.2; A, B, C), headings and footers. Your outline will print out per your specification with automatic page numbering. All text and graphics documents associated with your outline entries will be printed in the appropriate place in the outline. Unfortunately, all text is displayed in the same font, size, and style. A Text Settings option allows you to set these parameters but all text in the outline must be the same.

An outline by its very nature, is a big picture look at a subject. While there may be pages of detailed information backing up the outline, it is hidden. Outlook allows you to view as much or as little detail as you want to see. By double clicking any bullet, the next lower tier entries are hidden. Double clicking the bullet again will make them reappear. Triple clicking performs the same function for all entries under the selected entry regardless of their relative level. This feature looks smoothly and quickly to let you toggle between looking at the "forest" and the "trees."

Outlook works just as well with graphics as it does with text. You can place graphics under any heading and quickly view it full-sized by double clicking the related ballet. This feature could easily be used to construct an indexed library of clip art, charts, or other graphics.

## Let's Make a Plan

The latest version of SideKick also contains MacPlan, an effective little spreadsheet and graphics program in a desk accessory. MacPlan supports 50 rows by 20 columns. With up to 69 cells used for labels, a maximum of 931 cells are left for entering data. You may request MacPlan to format data as dollars, dollars and cents, percent, or as a number.

Use of this spreadsheet is straight forward. Any cell can be selected with the mouse or the cursor keys. Data may be entered directly or instead a formula may be added which can be used by MacPlan to calculate the value for a cell. For example, if a cell (say D7) were to contain "net profit," then you might indicate that C7 = C6 (gross sales) - C5 (expenses). Cell C6 (gross sales) might be calculated by MacPlan as the total of cells B1 through B12 (monthly sales). Then, if you revised your projection of July sales, the total sales and net profit figures would be automatically updated.

A wide range of formulas are provided for calculating the value of cells based on the value of other cells. In addition to basic operators (e.g., +, -, /,x, %), financial functions, depreciation functions, mathematical functions, logical functions, and calendar functions are included.

MacPlan allows you to create pie charts, line charts, and bar charts. Creation of a chart simply requires selecting "Graph" from the MacPlan menu, entering the title and legends, and indicating the type of chart you want. Bar charts may be plain, stacked, or grouped. You have no flexibility in determining fill patterns, fonts, on the size of the graph. Only 18 or so data points can be shown in pie charts and regular bar charts. With stacked or grouped bar charts up to 3 data entries can be made for each of the 18 points. For charts with many data points, part of the legend may be cut off and will not print. This feature has only limited usefulness

To print MacPlan spreadsheets and graphics you must copy the document then paste it into an application with printing capabilities.

## Just Keeping Time

I get tired of all those boring little digital clocks for the Macintosh. I know they are compact and easy to read, but they are not very interesting. I was glad to see that SideKick's timekeeper is a 1" x 2" analog clock with a second hand. The second hand continues to sweep even when another window is active. This clock seems to do what it was designed to do. It does not contain an alarm function.

### Monday, Monday

At first glance, the SideKick calendar looks similar to the simple public domain calendars available for the Mac. A closer look, however, indicates that it is much more. Scroll bars allow any date between 1905 and 2030 to be viewed. Notes may be associated with any date. Dates for which notes are added appear circled. The notes window scrolls vertically to allow entry of more text than can be continued in the window. Pressing the return key automatically prints a dashed line to separate different notes.

A potentially useful feature of SideKick's calendar is the "Find" function. Unfortunately the usefulness is only *potential* and not real. First, it only searches the notes of the date currently selected date. It would have been useful to search all dates in a certain date range to see when the air conditioner filter was last changed. Also, I tried a number of searches on the notes of the currently selected date and none of them were completed successfully.

The Calendar Menu allows you to select the "Week at a Peek" option. This option allows you to view the notes for 7 days at once. The menu also allows either the regular calendar window or the "Week at a Peek" window with associated notes to be printed.

### When You Need To Do Some Figuring

SideKick's calculator called "Calculator +" is a multifunction business calculator with "paper tape" output to the screen and printer. It allows entries from the keyboard or by clicking the keys on the calculator display.

Calculator + includes roots, logs, and trigometric functions as well as the usual mathematical functions. Also, it includes some very helpful financial function including:

- Calculate the growth of an initial investment
- Calculate the worth of a future return
- Calculate the worth of a future stream level of payment
- Calculate the level payment necessar to pay off a loan

Each of the above calculations can be made with only 2 or 3 entries.

Almost every Calculator + function can be operated from the keyboard which is much faster than using the mouse.

Calculator + is reasonably easy to read and use. It provides a new and powerful alternative to the Apple calculator.

### Area Code Lookup

Area Code Lookup is a small desk accessory that gives information about any area code entered into a dialog box. For example, when "205" is entered the window reports:

Locality: Alabama
Region: Birmingham
Time: Central

This may occasionally be useful in predicting whether a long distance contact may be a lunch or not.

### Write It Down

"Notepad +" is a mini word processor that creates text files up to 28k which are compatible with MacWrite and Microsoft Word. The Notepad + menu allows you to open existing text files, save a file, print a file, or find characters, words, or phrases in a file.

The Notepad + window can be sized for convenient use with other windows on the screen. This DA is designed to be closely integrated with other SideKick functions. For example, Notepad + files can be sent by modem with MacTerm and printed in the background using Readiprinter.

Note that since Notepad + works only with text files, information related to page format, fonts, and styles is not included.

### Send It By Wire

MacTerm is SideKick's Hayes-compatible 300, 1200, and 2400 band telecommunication desk accessory. It may be used to transmit and receive data while running other applications.

After selecting the "Configure MacTerm" icon, the usual adjustments may be made to match the software with your modem. Additionally, you may choose either the printer or modem port.

MacTerm remembers up to seven numbers that you frequently dial via modem and can automatically dial them for you.

### You Rang?

One of the best known SideKick components is its telecon management tool also calles SideKick. (There is an desk accessory called MacDialer that performs many of the same functions). This feature can help you manage and organize every part of your telephone communications.

When SideKick is selected, the screen is filled with an attractive display of information and command buttons. It allows you to:

- Enter a phone number to be dialed by click ing on appropriate numbers in the key board display. Alternately you may select a name from the scrolling "PhoneBook" win dow and have SideKick look up the number and dial it for you. Dialing may be via modem or PhoneLink™. You may set the di aling speed as required.

- SideKick will time your call and calculate phone costs (for long distance calls) as well as consulting fees.

- You may make notes on each call by typing into the "PhoneNotes" window. These notes are saved as part of your phone log and may print out with other information about calls.

- You may go from this screen to the SideKick calendar by clicking the calendar button.

- You have the ability to easily add, delete, or change entries in the phone book. At your option, phone book entries may also appear in a menu of names. Entries in the Phone Book may be sorted in a variety of ways. A "miscellaneous notes" space allows you to enter information about each entry (e.g. birthday, shirt size, product). Also, you may have a number of different Phone-Books, one for home, one for business con tracts, etc. You can then switch between the phone books as needed.

As you make and receive phone calls, SideKick maintains a log of each call including your notes. This log can be viewed, sorted in a variety of ways, and printed.

The SideKick telephone organizer function does even more than I have summarized above. It's not hard to understand why it is one of the most popular such programs on the market.

### QuikSheets

QuikSheets is a combination of four input forms that may be used to capture specific infor-mation. The forms included are:

- Alarms—Lets you enter the times of events you wish to be reminded of. When the alarm goes off a bell sounds and the Apple menu bar be gins to blink. It continues to blink until you acknowledge the alarm.

- Expenses—Provides for the accumulation of expenses with totals automatically calcu lated.

- Credit Cards—Provides for the recording of critical information about your credit cards.

- Things-to-do—Lets you enter things to do and check off accomplished items.

The QuikSheets desk accessory allows you to compact the cards by removing blank lines, com-pleted tasks, and past alarms. QuikSheets may be printed.

### More, More, More

In addition to the applications described above, SideKick contains a number of utilities. These utili-ties enhance the usefulness of the overall SideKick package. Included are:

- PrintManager—Allows you to obtain disired information from your Phone Book, CalendarBook, and Quik Sheet files and print it in a variety of formats.

- Converter—Allows you to copy phone book entries from Habadex, MacPhone, and various text files into your Phone Book.

- QuikEditor—Lets you design your own QuikSheets to keep track of all kinds of data.

### In Summary

These seem to be the days of less for more. Much of todays software doesn't deliver on all its promises yet prices continue to escalate. SideKick is a delightful exception. It delivers! Each function that I used did its job in a simple and effective manner. Except for a minor complaint about the graphing function, I found SideKick to be an excel-lent value. When you can get a program with all these functions for $59, you should grab it!

**EOF**

*FOR THOSE WHO NEED TO KNOW*    **68 MICRO JOURNAL**

# SOFTWARE

## USER

**A Tutorial Series**
By : Ronald W Anderson
3540 Sturbridge Court
Ann Arbor, MI 48105

## NOTES

*From Basic Assembler to HLL's*

## Assembler is Assembler is Assembler... or Not?

Last Saturday an envelope arrived in my mailbox. It was a copy of part of a letter received by '68' Micro Journal from someone in Freedom PA (signature illegible in the copy). I quote one paragraph from a section headed "A few comments". "Another idea that would be nice. How about printing a simple program in XBASIC, BASIC09, C, PL9, Flex Assembler, OS9 Assembler and Pascal. ... The program should print text, loop with conditions and (do) simple arithmetic." My first reaction was a mental objection to the separation and distinction between OS9 Assembler and FLEX Assembler. After I thought about it a little, the objection went away. The programs may be quite different, and must be somewhat different because of operating system requirements.

Unfortunately, I don't have "all of the above," but I thought I'd take on the challenge with what I have. First, I do have TSC Extended BASIC, so I thought I'd dream up a program and then write it in as many languages as I could. I came up with a program that prints the heading "Multiplication Table," a blank line, and then the table for 9 times 0 through 12. (See listing 1). Each line of the table is to read out "3 X 9 = 27" etc. The BASIC program is all of 6 small lines, including the END statement which is not really necessary. It is certainly at least one of the shortest source listings in the group. Remember that XBASIC is an interpreter and that the interpreter is about 16K of program. Rather than just show program listings I though it might

be informative to indicate how big the resulting programs are in terms of object code produced.

Lucidata Pascal is responsible for the next listing, (listing 2). just 7 lines of code. This Pascal is what is called a P-Code Pascal. It generates a code file that is sort of "semi-compiled." That is, the code is not machine code, but a sort of Pseudo code that may be interpreted by a simple interpreter. Because of this, the size of the object code is sort of indeterminate too. It is possible to compile and "bind" some of the runtime code to the compiled P-Code to produce a runnable object file, but a great deal of interpreter is included whether or not it is used, so the object code is larger than necessary. Please note that these two paragraphs are in no way a criticism of BASIC or P-Code Pascal, just an explanation of why object code size wouldn't mean much if reported here. Lucidata Pascal is a FLEX running 6809 compiler. I would hope that they might come up with a 68000 running version someday. Perhaps some of the following might serve to convince them to do so.

The programs are presented more or less in order of byte count. First let me say that unfortunately I have never run OS9 on the 6809 system. Perhaps some reader might fill us in on that version. I did write the program in 6809 Assembler for FLEX however. Listing 3 is that program. It uses three of the FLEX routines and this, of course is what would make the FLEX and OS9 versions different as you will

see below. In order to help you understand the program, I've inserted the BASIC program lines as comments that precede the section of code that implements that line. In one case I've split a BASIC line into three parts. In order to keep the code closer. I've written the program in position independent code. The program is about one page and the byte count is 116.

I could not resist writing the program in 68000 Assembler, since I have both SK*DOS (remarkably like FLEX in operation and user available routines) and OS9 for the 68008 machine. The SK*DOS version is just about a direct translation from the FLEX 6809 version except that there are so many data registers available that I didn't have to define any variables in the program. All operations are carried out using data registers. I note that I couldn't use the PSTRNG routine of FLEX because it puts a CRLF BEFORE the text string that it prints out. SK*DOS has a routine with the identical function but it also has a PNSTRN function that doesn't put in a CRLF. I therefore didn't have to code a PSTRNG routine as I did in the FLEX version. Note that the PNSTRN SK*DOS routine requires A0 to be pointing at the start of the string and that the string must end with $04. OUT5D translates a 16 bit binary integer number to a decimal ASCII string and outputs it. The binary number must be in D4 and D5 is set to zero to inhibit leading spaces. The resulting program is only 82 bytes of

object code, and is the champ in terms of brevity. I should mention that SK*DOS itself uses and can alter the registers A4-A7 and D4-D7. The other half, D0-D3 and A0-A3 are available to the user for his program. (See listing 4).

I had been toying with the idea of figuring out how to write assembler programs for OS9 use now that I had written some for SK*DOS, so I took it as a challenge one evening to write a program to print "Hi There" to my terminal. I had to read or at least skim about 100 pages of manual to find the necessary information, and between that, what seemingly was a sample program that wouldn't assemble properly, and a number of dumb errors on my part, it took about 4 hours to reach my goal. I was a bit frustrated but pleased. A day or two later I wrote the MULT program in OS9 68000 assembler. See listing 5.

The string of labels that follow the word "psect" are those necessary to generate the proper OS9 "header" for the file. The values for Edition, Typ_Lang and Attr_Rev are all defined in one of the files that is included by the "use defsfile". Psect stands for "Program Section", the place where non-alterable information is stored, that is, the program code and any constants may be put there. Vsect is the Variable section where alterable memory locations reside, that is locations where variables may be read or written.

OS9 is a multi-user system, and it has a nice feature. Two people can use the same copy of the program in memory at the same time because when they call it, they are each assigned a different variable area in memory. The register A6 contains the "base address" of the variables. When the computer switches tasks (users) it saves all the registers for the user whose task is temporarily suspended, and loads the registers with the values for the current task. A6 is different for each user of the same program. The variables are accessed by using the name as an offset from A6. (e.g. move.l d0,varname(a6)). This program has only one variable, called buffr. It is 14

bytes long and a reference to buffr(a6) will get the contents of the first byte, word or long depending on the reference. In the present case, A0 was set to point at the buffer by "lea buffr(A6),A0". Then A0 was manipulated to point to bytes in the buffer as it was loaded with characters.

One advantage of OS9 is that all the registers are available to the user program except A7 which is the stack pointer, and A6 which contains the pointer to the variable area. All the other registers may be used by the user program.

A disadvantage is that it places rather rigid constraints on the assembler programmer with regard to variables and accessing them. However, it is very good programming practice anyway, to separate the variable section from the program section. It can sometimes avoid accidentally writing over program instructions. OS9 guards data areas very carefully and any attempt to write to memory outside of your data area will cause an instant error message and the stopping of your program's execution.

The OS9 Assembler manual mentions that a file called sys.I should be used to include the system routine call names and values in the program. The manual wasn't specific about how to use the file, but it became apparent that it is a so called relocatable object file and that it is to be included at link time. The assembler produces relocatable object code, which must be linked by the linker before it can become an executable command file.

I decided to use the I$Write system routine which is similar to the PNSTRN call of SK*DOS. D0 must contain the "path number" for the output (I to output to the STDOUT terminal). D1 must contain the length of the string, and A0 must point at the first byte of it. The system function F$Exit will report an error whose number is loaded in D1 before exiting the program. D1 must therefore be cleared before calling F$Exit. The syntax OS-9 F$Exit is an assembler convenience.

The list of equates in the defsfile is scanned and the proper code placed in the file.

I wanted to use the math function that translates a long integer to an ASCII string, similar to OUT5H in SK*DOS, but I found that the routine includes all leading zeros, and the numbers output would have looked like "0000000018". I decided it would be simple to write my own conversion subroutine called "printi" which came out to 17 lines of code. What I have just said could be repeated for the "crlf" subroutine at the end of the program. All the register stuffing prior to the OS-9 call takes its toll in terms of program byte count, but remember that OS-9 doesn't tie up half of the registers for itself, and that it can do numerous tasks simultaneously. At any rate, the finished program was 174 bytes worth.

Now for the "native code" compiler versions. Long time readers might remember Whimsical, a compiler for the 6809 written by John Spray in New Zealand. Listing 6 shows the program, and again it is only a few lines. John has recently sent me a preliminary version of Whimsical to run under SK*DOS in 68000. The code required for this program is identical in both systems. The compiler produced 245 bytes of 68000 code and the 6809 version produced 328 bytes of 6809 code. Note that the 68000 version again was smaller than the 6809, just as happened in the FLEX / SK*DOS assembler cases. Though the 68000 instructions are more bytes, each one does more, and the same program can be done with less overall code. There are several ways to specify CR and LF in Whimsical, one is via ^M^J sequence. Another is CHAR $0D, CHAR $0A. Note that the Whimsical Syntax is pretty much in agreement with Pascal. John Spray has told me that this preliminary 68K version produces excess code and that a lot of optimization can be done to improve the object code size. I wonder how much better John can do?

Listing 7 is the PL9 6809 version and the OS9 PLuS 68000 version. The

same input code will work for both compilers again. About the only disadvantage to PLuS / PL9 is that print statements can't have a strung out list of arguments. A different print procedure must be called for strings than for numbers. It is possible in this language to trim the library routines (IOSUBS in the present case) to remove any procedures that the program doesn't use. I did that for both versions. The 6809 version is 365 bytes and the 68000 version about 1300. The OS9 interface would seem to be responsible for the larger code in that version. Remember, in comparing, that the OS9 system has the advantage of multi users and multiple tasks running at the same time. When a computer has memory in the area of 750K or so, a few hundred bytes in a program are not an extravagance.

Listing 8 is one of two "C" versions that I did. This one would seem to win hands down for the shortest source code, but at least in my opinion, it comes in dead last in the compilers for readability. "C" has major advantages, the most important one being that it is very "portable." Most of the features of "C" are contained in what is called the "Standard Library." This library is written in "C" so that it moves from computer to computer without major rewrites. However, because it is written in "C" and not assembler, it is a little bulky. This version of the MULT program in Microware OS9 "C" for the 68000 produced 11,300 bytes of output code. The 6809 version (identical source code) in Windrush "C" by James McCosh produced 5246 bytes.

I decided that the "printf()" library function must be responsible for a lot of the code, so I did another version using the fputs() function (fILE putsTRING) for the strings, directing the output to the terminal. I wrote a simple printint() function (which is recursive) patterned after (read that "stolen from") the routine by Graham Trott of the same name in IOSUBS in PL9. Anyway the program compiled to 5140 bytes in the 68K version (less than half the size of the original), and

to 4059 bytes in the 6809 version, about 25% smaller. See listing 9.

All of the programs produce essentially identical output with a few variations in number format and spacing. The heading was printed in keeping with the style of the language. Pascal and PL9 generally use upper case (no rules about it, just tradition, at least for me) so the headings are all caps. "C" programmers for some obscure (to me) reason prefer lower case as a convention for program statements and all caps for defined constants. The heading is all lower case in that program.

Generally, a program in assembler tends to be larger in source code and smaller in object code. It takes longer to debug an assembler program simply because there are more instructions to consider as a source of error. There are also more opportunities to make errors because more instructions must be used. The higher the level of the language, the smaller the source code. Some languages tend to be "self documenting," for example, the Whimsical, PL9 and Pascal versions. Though someone who has spent the past 5 years writing "C" programs every day would disagree, "C" programs tend to be more cryptic to the reader. Assembler programs may be done clearly if written in small "chunks." I see no reason to go as far as a main program that does nothing but call subroutines when numerous subroutines will only be used once, but it is a good idea to separate or group in-line code by function with a comment line or two between the various portions. Of course if a particular function is needed more than two or three times in a program a subroutine is called for simply to reduce the code produced. A good side effect is to improve the readability of the code.

Now a word about recursive procedures. Note that the printint() in the "C" program references itself. Most of the languages that are considered "procedural" in nature (they break a program into functions or procedures or subroutines) allow two kinds of

variables. There are "Global" variables and there are "Local" ones. Global variables are "declared" at the beginning of the program and they exist as long as the program is running. Local variables are associated with particular procedures or subroutines. When the procedure is running, space for local variables is allocated (usually on the stack). When the procedure ends, the space is deallocated. Local variables won't retain information once the procedure is finished and control returned to the main program ("c" has a special case called a "static" variable that may be defined within a function, that remains between calls to that function). At any rate, this method of allocating variable space for a procedure when it is "called" leads to an interesting possibility. In the present case, look at the printint() function in listing 9. The function is passed the number to be converted from binary to ASCII characters for printing. The function has a local variable called "digit." digit = number % 10 + ox30; This statement (% is the "mod" or modulus function. Number Mod 10 results in the REMAINDER resulting from dividing number by 10, which is the lowest order digit of the number to be printed. We add $30 to it and the result is the ASCII code for that digit. Then we divide number by 10 (integer divide) which essentially throws away the remainder that we just turned into a character. That character is placed in the local variable "digit" remember. ,p Now comes the tricky part. If the number is not zero, (number, which was passed to the function is also a local variable on the stack), the function calls itself, creating new local variables number and digit in a new place on the stack. The value of number passed to this second call to printint() is not the starting value but the value resulting from the divide by 10 in the first call to printint(). The new call to the procedure again does the mod function, placing the next higher order digit in the new local variable "digit." The key to what is going on is to remember that the first call to the function printint has not yet finished executing. Its execution was interrupted by the second call from

within the function itself. The process continues, each call to printint() getting as far as the next call, number getting smaller each time by the successive divides by 10, until number becomes zero. At that point the last call to printint() goes on and executes the last line, putc (digit, stdout);. As soon as that happens, the last call to printint() is finished and it returns to the previous call which has the next digit in its local variable "digit" It finishes by printing the character and returns to the previous call until the first call is reached when it prints the last character and returns to the main program.

Note that the OS9 68K assembler version of this routine does essentially the same thing but it is not recursive. It puts the digit that it has calculated at the end of a buffer and decrements a pointer, then continues in a loop to calculate earlier digits until number is reduced to zero. Then it loads the digit count and calls the print string function of OS9. One could make the assembler version recursive simply by pushing the currently calculated digit onto the system stack move.b dn,-(a7). Because of the way a character must be output with I$Write, when the last call to the procedure is completed, one would have to pull the characters off the stack and stuff them in a buffer to have them assembled for printing anyway, and that would just be more code.

If all this is thoroughly confusing remember that "C" has the built-in function to print a number but we were trying to be a little more efficient and use our own. Pascal, Whimsical, and PL/9 all have such functions as well.

I hope this discussion has given you a little glimpse of programming from several points of view. I point out that several of the programs are identical in 6809 and 68000 versions when written in one of the higher level languages. The assembler code is quite different for the two processors, and even with one processor, the code is significantly different depending on the operating system. Assembler code is very efficient, runs fast, uses less bytes of code, but it is also VERY specific to the processor and operating system. The higher level languages shield the user from the processor but not generally quite completely from the operating system. The format and allowable size of filenames, for example are usually a function of the operating system and not of the language. Some of the languages allow you to do a "system call," which of course is operating system dependent. For example XBASIC will let you EXEC,"ITYSET PS=N" to turn the pause function off before outputting to a printer. Microware "C" allows system calls to OS9 to do things like rename or delete a file or to set up terminal parameters. The call system "nopause" does the same thing as the EXEC example for XBASIC above.

**Listing 1**

```
10 PRINT "MULTIPLICATION TABLE"
12 PRINT
15 FOR N=0 TO 12
20 PRINT N;"X 9 =";
30 PRINT N*9
40 NEXT N
```

**Listing 2**

```
PROGRAM MULT (OUTPUT);

    VAR NUM, RESULT:INTEGER;

BEGIN
    WRITELN('MULTIPLICATION TABLE');
    WRITELN;
    FOR NUM:= 0 TO 12 DO WRITELN(NUM,' X 9
-',NUM*9);
    END.
```

**Listing 3**

```
 NAM MULT
 TTL PROG TO PRINT MULTIPLICATION TABLE
*
* FLEX EQUATES
*
OUTDEC EQU $CD39
PCRLF  EQU $CD24
PUTCHR EQU $CD18
WARMS  EQU $CD03
```

```
 ORG 0
START BRA BEGIN
MSG FCC /MULTIPLICATION TABLE/,$04
MSG1 FCC / X 9 = /,$04
NUM RMB 2 SAVE TWO BYTES FOR N
RESULT RMB 2
* 10 PRINT "MULTIPLICATION TABLE"
BEGIN LEAX MSG,PCR
 BSR PSTRNG
 JSR PCRLF
* 20 PRINT
 JSR PCRLF
* 30 FOR N=0 TO 12
 CLR NUM,PCR
 CLR NUM+1,PCR
LOOP LDD NUM,PCR
 CMPB #12
 BGT ENDLOP
*40 PRINT N;
 LEAX NUM,PCR
 CLRB
 JSR OUTDEC
* 40 CONTINUED PRINT"X 9 = ";
 LEAX MSG1,PCR
 BSR PSTRNG
* 40 CONTINUED PRINT N*9
 LDD NUM,PCR
 LDA #9
 MUL
 STD RESULT,PCR
 CLRB
 LEAX RESULT,PCR
 JSR OUTDEC
 JSR PCRLF
```

```
   *50 NEXT N
    LDD NUM,PCR
    ADDD #1
    STD NUM,PCR
    BRA LOOP
ENDLOP JMP WARMS
*
* PSTRNG SUBROUTINE
*
PSTRNG LDA ,X+
    CMPA #$04
    BEQ DONEP
    JSR PUTCHR
    BRA PSTRNG
DONEP RTS
*
    END START
```

### Listing 4

```
    NAM MULT
    TTL MULTIPLY
*
* SK*DOS EQUATES
*
PNSTRN EQU $A036   A4 IS POINTER
OUT5D EQU $A038    D4 WORD D5=0 TO SUPPRESS
SPACES
PCRLF EQU $A034
WARMST EQU $A01E
*
    ORG 0
START BRA BEGIN
*
MSG DC.B "MULTIPLICATION TABLE",$04
MSG1 DC.B " X 9 = ",$04
* NUM WILL BE IN D0
* RESULT WILL USE D4
*
* 10 PRINT "MULTIPLICATION TABLE"
BEGIN LEA MSG(PC),A4
    DC PNSTRN
    DC PCRLF
* 20 PRINT
    DC PCRLF
* 30 FOR N=0 TO 12
    CLR D0  N
LOOP CMP #12,D0
    BGT.S ENDLOP
* 40 PRINT N;
    MOVE D0,D4
    CLR D5
    DC OUT5D
* 40 PRINT "X 9 ="
    LEA MSG1(PC),A4
    DC PNSTRN
* 40 PRINT N*9
    MOVE #9,D4
    MULU D0,D4
    CLR D5
    DC OUT5D
    DC PCRLF
* 50 NEXT N
```

```
    ADD #1,D0
    BRA.S LOOP
ENDLOP DC WARMST
    END START
```

### Listing 5

```
    nam mult
    ttl os9 mult prog
    use defsfile

    psect mult,Typ_Lang,Attr_Rev,Edition,0,Mult

    Edition equ 1
    Typ_Lang equ (Prgrm<<8)+Objct
    Attr_Rev equ (ReEnt<<8)+0

    StdOut equ 1

    message dc.b $0d,$0a,"Multiplication
table",$0a,$0a,$0d
    mesglen equ *-message
    mesg1   dc.b " X 9 = "
    msg1len equ *-mesg1

    vsect
buffr   ds.b 14
    ends

* 10 print "Multiplication Table"
* 20 print
Mult    moveq.1 #1,d0  path number
        move.1 #mesglen,d1 number of bytes
        lea message(pc),a0  buffer pointer
        os9 I$Write
* 30 for n=0 to 12
        clr d2
Loop    cmp #12,d2
        bgt.s Endlop
* 40 print n," X 9 = ",n*9
        move.1 d2,d1 pass number in d1
        bsr.s printi print an integer
        moveq.1 #1,d0
        move.1 #msg1len,d1
        lea mesg1(pc),a0
        os9 I$Write
        move.1 d2,d0
        move.1 #9,d1
        mulu.w d0,d1 product in d1
        bsr.s printi
        bsr.s Crlf
        addq.w #1,d2
        bra.s Loop
Endlop  move.1 #0,d1
        os9 F$Exit

* Enter with number in d1 convert and put in
buffr
* and print via I$WritLn
*
printi  lea buffr(a6),a0
        adda.1 #12,a0
        clr.1 d3 counter for bytes in string
```

```
ploop    move.l #10,d0
         divu d0,d1 mod in upper word of D1
         swap.w d1
         add.b #$30,d1
         move.b d1,-(a0)
         clr.w d1 clear remainder
         addq #1,d3
         swap.w d1
         tst.w d1
         bne.s ploop
         move.l #1,d0 path
         move.l d3,d1 length of string
         os9 I$Write
         rts
* crlf routine
Crlf     move.w #$0d0a,buffr(a6)
         lea buffr(a6),a0
         moveq.l #1,d0
         moveq.l #2,d1
         os9 I$Write
         rts

         ends
```

**Listing 6**

```
% mult program in whimsical 6809

BEGIN
  INTEGER NUM;

  WRITE "MULTIPLICATION TABLE^M^J^J";
  NUM :=0;
  WHILE NUM <=12 DO BEGIN
    WRITE NUM," X 9 = ",NUM*9,"^M^J";
    NUM :=NUM+1;
  END;
END.
```

**Listing 7**

```
/* PROGRAM MULT IN PL9 */

ORIGIN = 0;
STACK = $BE00

GLOBAL INTEGER NUM;

INCLUDE IOSUBS.LIB.1;

PROCEDURE MAIN;
   PRINT "MULTIPLICATION TABLE";
   CRLF;
   CRLF;
```

```
   NUM=0;
   WHILE NUM < 13 BEGIN
      PRINTINT(NUM);
      PRINT " X 9 = ";
      PRINTINT(NUM*9);
      CRLF;
      NUM = NUM+1;
   END;
```

**Listing 8**

```
/* mult table program in C */


#include <stdio.h>

main()
{
    int num;

    printf("multiplication table\n\n");
    for (num=0; num <=12; num++) printf("%2d
X 9 = %3d\n",num,num*9);
}
```

**Listing 9**

```
/* mult table program in C 6809 */

#include <stdio.h>

main()
{
    int num;

    fputs("multiplication table\n\n",stdout);
    for (num=0; num <=12; num++)
    {
        printint(num);
        fputs(" X 9 = ",stdout);
        printint(num*9);
        putc('\n',stdout);
    }
}

printint(number)
    int number;
{
    char digit;
    digit = number %10 + '0';
    number = number/10;
    if(number !=0) printint(number); /* recur-
sive */
    putc(digit,stdout);
}

EOF
```

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# FORTH

## THE DRAG STRIP

I know that some of you must think that I have a fixation on execution speed, but I have found the topic to be fascinating. There are very practical reasons to be concerned about execution speed in FORTH, since it is often chosen as the embedded language for controllers because of its high speed. Furthermore, we are all interested in having our applications run as fast as possible, because we are irritated beyond all reason by having to wait on a computer. Therefore, I have continued my study of ways to get from here to there as fast as possible, hence the subtitle "the drag strip."

In this installment, I want to discuss ways of speeding up multiplication operations. Examination of Figure 8 in my last column shows that the word 2* executes about 6.6 times faster than the logically equivalent pair of words, 2 * . This difference cries out for investigation!

This is typical of a situation where a generalized technique is compared to a specialized case. In other words, 2* consists only of load, add, and store operations, so it must be faster than the much more complex, and thereby general, * operation.

### High Level FORTH

This set of experiments had two goals: (1) to speed up simple multiplication operations, and (2) to do it in high-level FORTH. Goal #1 is obvious, but goal #2 was chosen to maintain portability, even from processor to processor. In other words, I wanted to adhere to the principle of KISS (Keep it simple, stupid!).

Any multiplier greater than 2 can be generated by a series of doublings and algebraic additions. For example, consider 13 as a multiplier, as in Figure 1.

If you follow through the FORTH logic, you will see that the FORTH notation is identical to the logic of the last line of the algebraic notation. Any other multiplier can be generated in the same way. The FORTH screens 20 and 22-24 show this.

Screen 21 shows the actual timing loop used for testing 10* .

The multiplicand, 23, was defined as a constant by the line 23 CONSTANT 23 . There is no particular significance to using 23 as opposed to some other value; I just happened to choose it at random. The testing was done on my CoCo3 at 1.89 MHz. The external timer was controlled by the cassette recorder relay through the words TIME_ON and TIME_OFF . The data do show a bit more scatter than I would like, but the total time must be accurate to within +/- 0.5 second. I decided not to pursue this scatter any further at this time, since the data are certainly good enough to prove the point I am trying to make. The tests and results are summarized in Figure 2.

Notice that the definition of 15* differs from the others in that it contains a subtraction, rather than an addition. I did this simply in order to shorten the definition. The definition would have needed two more terms in order to contain an addition, instead of a subtraction, and this would

have also lengthened the execution time.

The last line of Figure 2 contains a measurement of the time required to multiply by 15 as one would normally do it. Notice that this operation was about 18 microseconds shorter per pass through the loop than the loop containing 15* . This simply proves that one can outsmart himself if he is not careful! If the definition gets too long, about 7 terms, then it is quicker to run the conventional path. However, notice that this definition of 16* is about 42 microseconds faster than 15 * , so it obviously pays to look for special cases.

### FORTH plus assembly language

Now let us consider the case where speed is most important, and we don't care about transportability. To do this, we should write the critical part of the definition in assembly language. This should give us the minimum execution time, short of writing the whole application in assembly.

This also demonstrates a great virtue of FORTH, in that it is so

very easy to incorporate assembly language definitions into a program. Sure, you can do the same kind of thing with C or Pascal, but not nearly so easily. The important points to observe here are that the assembly language definition is called just like any other definition and data are transferred on the Data Stack just the same way as with high-level definitions.

Since I don't have access to a 68000 system, I will have to stick to the 6809. Fortunately, FF9 has an assembler, so that part of the job will be easy. I am not a whiz at writing math routines in assembly language, so my efforts may not be the most efficient, but they do work. If anyone has faster routines, please share them with us.

The definitions I wrote are shown in screens 26-38. With the exception of the definition for 15* , all of the definitions are direct transliterations from the high-level definitions. I simply took Wilson Federici's FF9 definitions for DUP , + , and 2* and incorporated the important parts into each new definition. The definition of 15* now follows the pattern of the rest of the definitions as ( ( 7 * 2 ) + 1 ).

The test results are summarized in Figure 3.

These definitions average about 2.4 times faster than the corre-

sponding high-level definition, even the very long 15* . I don't think that you could make these operations run much faster, within the constraints of calling data from a stack and calling the multiplier as a virtual subroutine. I think that in-line code would be the only way to go any faster!

### Caveats

Caution, this is signed integer math, limited to 16 bits. In other words, the product of any multiplication can never fall outside the range of +32767 through -32768. Any product exceeding these limits will, at best, be truncated on the high end, and, at worst, bear no resemblance to any part of the real answer. Remember that FORTH has a different set of definitions for use with 32- bit numbers.

Also, these definitions can only be used where the multiplier is a constant. If you must multiply two variables, you must use the conventional * operation. High-level FORTH versus assembly language and others

The factor of 2.4 penalty paid for running in high-level FORTH is hardly any penalty at all!. What other high-level language can do so well? Of course, now that I have the assembly definitions written, I will probably use them most of the time, instead of the high-level versions, but I will

probably rarely, if ever, notice the difference. In the cases shown here, most of the added time comes from the overhead encountered in "calling subroutines." Since C and Pascal function, mainly, by also calling subroutines, the compiled form of either one could hardly be expected to be faster than well written FORTH. For what it is worth, I think that it is harder to write good, tight programs in C or Pascal than it is in FORTH, so that is how I account for the general pronouncement that FORTH is usually faster than either one. By the way, I have seen threaded FORTRAN execute faster than conventionally compiled FORTRAN, both on a PDP-11. Besides, FORTH is more fun!

### MORE ON FLEX AND THE COCO3

I forgot to mention in my last column that you must use the WIDTH32 screen option when you are loading FLEX into your CoCo3, if you have the same version as I use. If you try to load FLEX while in WIDTH40 or in WIDTH80, the system will lock up, and you will have to reboot. It took me a while to work this one out, so I can't imagine why I forgot to stress it last time. Once loaded, you can select the V51.CMD display, just as if you were using a CoCo1 or 2, but you cannot use the "hardware" 80 or 40

column screens. The problem is that FLEX clobbers some of the pointers and drivers when it loads; you must have a new driver if you want an 80 column screen. No, I have not written one!

I also said last time that I was having trouble getting my printer to work. Well, that problem has gone away, and I have no idea why!?! In any case, the printer works fine, now.

### FORTH WITH OS9, LEVEL I

I have now had a chance to spend some time on the DELPHI FORTH. It runs quite well on OS9, level I, version 2.0. This FORTH is called FORTH09 by the author, Dennis M. Weldy. He says that it is a preliminary version, which I hope means that he will soon publish a more complete version.

Once you figure out some of the peculiarities forced by OS9, this is a very "user friendly" FORTH. It has no editor directly associated with it. Developing programs is no real problem, since you can enter trial definitions directly from the keyboard. Once you have a definition working properly, you use the normal OS9 editor, or any other text editor, to prepare an ASCII file, such as the one shown in Figure 4.

As you would expect, this simple program prints the

message and the numbers 0 through 99. The indentation is not required, but makes the definition easier to read. Since this is a text file, there are no screen numbers to be concerned with, which some people consider an advantage (I think that the jury is still hung on that question). The program is called by the following command:

```
FORTH09 <TEST1
```

Notice that this is a standard OS9 command line with input redirected to the disk file named TEST1. That is why the last line of the file consists of the name of the definition to be executed; I don't know of another way to get input into the FORTH09, once you have redirected the input. In any case, this works nicely, and I have no complaints.

Since this is a "preliminary" version, and OS9 frowns on the practice, there are no words for direct interaction with the machine. Table 1 lists the words which are available, and you can sure do a lot with what is here. Certainly, if you are a good C programmer, you can add more words, since the source code is available.

The basis for FORTH09 was the first edition of Brodie's "Starting FORTH." Therefore, a few of the words are from Polyforth or fig-FORTH, and not FORTH- 83, but that is a minor quibble. The bigger problem is the lack of fetch and store definitions.

All in all, I like FORTH09 for what it is, but more for what it can become!

**EOF**

```
    Algebraic notation:
        13 = ( 2 * 6 ) + 1
        13 = ( 2 * ( 2 * 3 ) ) + 1
        13 = ( 2 * ( 2 * ( ( 2 * 1 ) + 1 ) ) ) + 1

    FORTH notation:
        13* = DUP  2*  DUP  2*  +  2*  +
```

*Figure 1. Derivation of 13* .*

| multiplier | | | | | | | | measured time(sec.) | adjusted time(micro) |
|---|---|---|---|---|---|---|---|---|---|
| 3* | ( DUP | 2* | + ) | | | | | 147.0 | 85.5 |
| 4* | ( 2* | 2* ) | | | | | | 131.3 | 69.8 |
| 5* | ( DUP | 2* | 2* | + ) | | | | 167.1 | 105.6 |
| 6* | ( 2* | DUP | 2* | + ) | | | | 167.0 | 105.5 |
| 7* | ( DUP | 2* | DUP | 2* | + | + ) | | 202.8 | 141.3 |
| 8* | ( 2* | 2* | 2* ) | | | | | 151.5 | 90.0 |
| 9* | ( DUP | 2* | 2* | 2* | + ) | | | 188.0 | 126.5 |
| 10* | ( DUP | 2* | 2* | + | 2* ) | | | 187.1 | 125.6 |
| 11* | ( DUP | DUP | 2* | 2* | + | 2* | + ) | 223.0 | 161.5 |
| 12* | ( 2* | DUP | 2* | + | 2* ) | | | 187.2 | 125.7 |
| 13* | ( DUP | 2* | DUP | 2* | + | 2* | + ) | 222.9 | 161.4 |
| 14* | ( DUP | 2* | DUP | 2* | + | + | 2* ) | 222.9 | 161.4 |
| 15* | ( DUP | 2* | 2* | 2* | 2* | SWAP | - ) | 230.8 | 169.3 |
| 16* | ( 2* | 2* | 2* | 2* ) | | | | 171.5 | 110.0 |
| 15 * | | | | | | | | 213.0 | 151.5 |

*Figure 2. Summary of the timing test results for high-level definitions.*

| multiplier | measured time(sec.) | adjusted time(micro) |
|---|---|---|
| 3* | 97.2 | 35.7 |
| 4* | 89.3 | 27.8 |
| 5* | 105.6 | 44.1 |
| 6* | 105.6 | 44.1 |
| 7* | 122.4 | 60.9 |
| 8* | 97.8 | 36.3 |
| 9* | 114.6 | 53.1 |
| 10* | 114.6 | 53.1 |
| 11* | 132.1 | 70.6 |
| 12* | 114.6 | 53.0 |
| 13* | 132.2 | 70.6 |
| 14* | 131.3 | 69.8 |
| 15* | 148.0 | 86.5 |
| 16* | 106.7 | 45.2 |

*Figure 3. Summary of the timing test results for assembly language definitions.*

```
    : TEST1   ( - )
              CR  ." This is a test of FORTH09 by D M
Weldy."  CR
              100  0  DO
                   I  .
              LOOP ;


        TEST1


        Figure 4.  Sample FORTH09 input.
```

```
      +         -         *         /         MOD
      /MOD      SWAP      DUP       OVER      ROT       DROP
      CR        SPACES    SPACE     EMIT      2SWAP     2DUP
      2OVER     2DROP     ."        =         <         >
      0=        0<        0>        NOT       AND       OR
      ?DUP      ABORT"    ?STACK    IF        ELSE      THEN
      I-        1-        2-        2+        2/
      ABS       NEGATE    MIN       MAX       >R        >R
      I         1'        J         */        */MOD     DO
      LOOP      +LOOP     LEAVE     BEGIN     UNTIL     WHILE
      REPEAT    U.R       PAGE      QUIT
```

Table 1.  The 64 primitive definitions within FORTH09.

```
SCR #20
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

  23 CONSTANT 23

  : 3* ( n - )                           \ RDL 02/11/88
    DUP 2* + ;

  : 3x ( n - )                           \ RDL 02/11/88
    3 * ;

SCR #21
   ( Timing loop                         \ RDL 02/10/88)

  : TEST  ( - )                          \ RDL 02/11/88
    TIME_ON
    20 0 DO
       $0000 0  DO
  23 11*
       SP!
       LOOP
    LOOP
    TIME_OFF
    7 EMIT ;

SCR #22
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

  23 CONSTANT 23

  : 4* ( n - )                           \ RDL 02/11/88
    2* 2* ;

  : 5* ( n - )                           \ RDL 02/11/88
    DUP 2* 2* + ;

  : 6* ( n - )                           \ RDL 02/11/88
    2* DUP 2* + ;

  : 7* ( n - )                           \ RDL 02/11/88
    DUP 2* DUP 2* + + ;

SCR #23
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

  : 8* ( n - )                           \ RDL 02/11/88
    2* 2* 2* ;

  : 9* ( n - )                           \ RDL 02/11/88
    DUP 2* 2* 2* + ;

  : 10*    ( n - )                       \ RDL 02/11/88
    DUP 2* 2* + 2* ;

  : 11*    ( n - )                       \ RDL 02/11/88
    DUP DUP 2* 2* + 2* + ;

  : 12*    ( n - )                       \ RDL 02/11/88
    2* DUP 2* + 2* ;

SCR #24
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88
```

```
  : 13*    ( n - )                       \ RDL 02/11/88
    DUP 2* DUP 2* + 2* + ;

  : 14*    ( n - )                       \ RDL 02/11/88
    DUP 2* DUP 2* + + 2* ;

  : 15*    ( n - )                       \ RDL 02/11/88
    DUP 2* 2* 2* 2* SWAP - ;

  : 16*    ( n - )                       \ RDL 02/11/88
    2* 2* 2* 2* ;

SCR #25
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

   \ DUP                                 \ MMF 02/11/88
   \    0 ,U LDD   _D PSHU

   \ +                                   \ MMF 02/11/88
   \    _D PULU  0 ,U  ADDD  0 ,U STD

   \ 2*                                  \ MMF 02/11/88
   \    0 ,U LDD  0 ,U ADDD  0 ,U STD

SCR #26
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

   CODE 3*       ( n1 - n2 )             \ RDL 02/11/88
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      NEXT,
   END-CODE

   CODE 4*       ( n1 - n2 )
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      NEXT,
   END-CODE

SCR #27
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

   CODE 5*       ( n1 - n2 )             \ RDL 02/11/88
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      NEXT,
   END-CODE

SCR #28
   \ Math experiments-MULTIPLICATION      \ RDL 02/11/88

   CODE 6*       ( n1 - n2 )             \ RDL 02/11/88
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      NEXT,
   END-CODE

SCR #29
   \ Math experiments-MULTIPLICATION      \ RDL 02/12/88

   CODE 7*       ( n1 - n2 )             \ RDL 02/12/88
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      NEXT,
   END-CODE

SCR #30
   \ Math experiments-MULTIPLICATION      \ RDL 02/12/88

   CODE 8*       ( n1 - n2 )             \ RDL 02/12/88
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD      \ 2*
      NEXT,
   END-CODE

SCR #31
   \ Math experiments-MULTIPLICATION      \ RDL 02/12/88

   CODE 9*       ( n1 - n2 )             \ RDL 02/12/88
      0 ,U LDD   _D PSHU                 \ DUP
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      0 ,U LDD   0 ,U ADDD  0 ,U STD     \ 2*
      _D PULU  0 ,U  ADDD  0 ,U STD      \ +
      NEXT,
   END-CODE

SCR #32
   \ Math experiments-MULTIPLICATION      \ RDL 02/12/88
```

```
CODE  10*     ( n1 - n2 )                                        \ RDL 02/12/88
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
    NEXT,
  END-CODE

SCR #33
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  11*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
    NEXT,
  END-CODE

SCR #34
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  12*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
    NEXT,
  END-CODE

SCR #35
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  13*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
    NEXT,
  END-CODE

SCR #36
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  14*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
    NEXT,
  END-CODE

SCR #37
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  15*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD   _D PSHU                              \ DUP
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      _D PULU  0 ,U ADDD  0 ,U STD                    \ +
    NEXT,
  END-CODE
```

```
SCR #38
  \ Math experiment-MULTIPLICATION                                \ RDL 02/12/88

  CODE  16*     ( n1 - n2 )                                       \ RDL 02/12/88
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
      0 ,U LDD  0 ,U ADDD  0 ,U STD                   \ 2*
    NEXT,
  END-CODE

+++
```

# BETTER WINCHESTER DRIVERS

By: Samuel I. Green, Ph.D.

In December 1986, the 68 Micro Journal published my paper called Simple Winchester which uses the Western Digital WD1000-05 controller. My article included modifications to drivers previously published by Graves and Zeff and was issued as Reader Service Disk-30. Since the article was published, I have changed drivers. I now use the Winchester drivers provided by Leo Taylor. In response to a letter in the 68 Micro Journal, I sent Leo Taylor six blank diskettes and he filled them with a treasure trove of utilities and documentation which I still haven't finished perusing. Included was a set of Winchester drivers far superior to mine. I changed over in 1986 and have never looked back.

By comparison, my drivers allocate one volume for each of two heads. WINSYS uses all available space on multiple drives and allows assignment to as many as 26 total volumes. My drivers never really worked right at 2 MHz even with I/O slowdown as required by the controller. WINSYS works fine at 2 MHz.

I finally contacted Leo Taylor about submitting these drivers to the 68 Micro Journal. He said I can submit them to you if he is credited as the source. I herewith submit Leo Taylor's Winchester drivers and a utility called SQUEEZE which is needed to modify FLEX for booting directly from the Winchester. I urge you to discontinue distributing my Reader Service Disk-30 and replace it with the enclosed diskette. *(editors's note: Done)*

WINSYS is a group of seven files to support Winchester drives connected to a WD1002-05 which has error detection and correction or to a WD1000-05 which does not. No change is required in the hardware interface described in my article. WINSYS includes the formatter, drivers, assigner, booter, parker, and documentation. Multiple Winchester drives are supported and a maximum of 26 volumes may be assigned. By appending the drivers to FLEX with the SQUEEZE command and adding a short boot to my ROM, I boot directly from the Winchester. All this was done by following the extensive documentation which is included.

I suggest that you print the .DOC files to get the complete description. Keep up the good work.

Sam Green
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*Editor's Note: Our thanks to Doctor Green and Leo Taylor for the following. They have been a couple of very knowledgeable contributors to the pages of 68 Micro Journal for many years. We appreciate their sharing their knowledge and programming "goodies" with all of us. Thanks Doc and Leo! We need good solid hardware and software articles like these.*

*There is no reason in the world why our S50 bus systems can't be upgraded to state-of-the-art systems. The 6809 has a lot of kick left. What I wonder is, "do we?" How about some more stuff of this quality from some of you other readers, please write or call for our writers guide, or even simpler, just go by the guide-lines on the content page of this issue. Remember, without your support we are nothing. Your contributions are what keeps 68 Micro Journal going!*

*DMW*

# WINCHESTER DRIVE SYSTEM SOFTWARE

By: Leo Taylor

## OVERVIEW

The Winchester System Software (called WINSYS) is a series of utilities for handling Winchester disk drives using the FLEX or STARDOS operating systems. WINSYS consists of the following:

WINDRV - the disk drivers WINFMT - the disk formatter WINASN - assigns Winchester volumes WINPARK - parks the head over landing zone WINBOOT - bootstrap loader for ROM monitor WINTABLE - user system equates WIN.DOC - this document WINEXAM.DOC - examples for this document

This group of programs provides several features. The system supports maximums of four drives, 1024 tracks per drive, 26 volumes, 256 tracks per volume, and 10 drive numbers. Each volume can be assigned to any drive number at any time using WINASN. FLEX can be bootstraped from a winchester drive when WINBOOT is added to EPROM. Minor features include proper control of the drive select LED, inter-volume protection, compact driver (about 300 bytes), command line parameters, readable source code, easy driver install.

The software was written to be easy to use, but since it supports many volumes one must be careful not to overwrite a volume. Do not try to format the drive until reading the section on setting up the table.

### !!! IMPORTANT NOTE !!!

If you edit WINTABLE.LIB you must reassemble BOTH the format and assign programs. If they get out of sync YOU WILL LOSE DATA.

## GLOSSARY

Drive number - FLEX drive number, usually 0-3. Physical drive - winchester or floppy drive select number, from a switch on the drive, usually 0-3. Logical volume - section of the Winchester designated by a volume letter, usually A-Z. The physical drives can have up to 26 volumes. A drive number can have any volume letter assigned. Drive LED - red indicator on the front of the drive. Maximum track - highest track used for data. Parking track - inside track, slightly higher than maximum track. Error correct - feature of some controllers that fixes data errors. EXAMPLE A - printout obtained from a simple system with one disk drive. Each part is numbered, EXAMPLE A1 is the WINTABLE.LIB listing, A2 is the WINFMT listing, and EXAMPLE A3 is the WINASN listing. EXAMPLE B - printout obtained from my system with two drives and 24 volume letters.

## HARDWARE SUPPORTED

Since I only know a few FLEX users with hard disks, I do not know the limits of WINSYS. The program should work with any hardware compatible with the following tested hardware:

1. Western Digital 1000-05 and 1002-05 controllers (not SASI). 2. Shugart 604 (5 meg) and Seagate 225 (20 meg) drives. 3. Three chip SS-30 interface (see schematic) and David Graves SS-50 board (68 MJ, Oct 82).

## UNDERSTANDING VOLUME LETTERS

The entire WINSYS is based on the concept of logical volumes which you must understand to use the programs. A logical volume is a section of the hard drive which may be from 2 tracks to 256 tracks. This program assigns a letter to each volume. In EXAMPLE A, letter B is a volume with 70 data tracks starting at track 31. Let's assume we formatted volume B with WINFMT and assigned it to drive number 1 with WINASN. If we catalog this drive (using CAT 1) FLEX will look for a directory on track zero of drive 1. Unknown to FLEX, WINDRV will add 31 to the track number (volume B starts at 31). The directory of volume B will be found at this track, and FLEX will display a catalog of the volume. If we save a file on this drive (using SAVE 1.TEST 100 200) FLEX will want to start the file on track 1 sector 1. WINDRV will add 31 to this track number also, so the file will actually start on track 32.

The important things to realize are:

1. The operation of WINDRV adding offsets to track numbers is transparent to FLEX. No FLEX program can detect the offset to the physical track. 2. Each FLEX drive number can only be assigned one drive letter (track offset) at one time. WINASN is used to change the letter (and track offset) assigned to a drive number. 3. Volumes can be as short as one data track (2 tracks if you count the directory) or as long as 255 data tracks. 4. Since a drive can have many volumes of up to 256 tracks, a drive of up to 1024 tracks (about 65 Megabytes) can be supported under FLEX.

## SETTING UP WINTABLE.LIB

The file WINTABLE.LIB is a library file used by most of the programs in WINSYS. It must be present on the work drive to assemble any of the other programs. It contains all of the user defined system equates, as well as the volume sizes. It MUST be edited before you can begin assembling the programs.

Before you assemble any of the programs, you must consider what size volumes you wish to have. It will take awhile (and perhaps several false starts) to get the sizes the way you need them. The volume sizes can be as small as one data track or as large as 255 data tracks. Though one track seems too small to be of value, bear in mind that with a 4 head drive one track is 128 sectors (over 1/3 of a SA-400 floppy). For example, I use a single track volume called LIBRARY for the 'C' compiler library files.

The first half of WINTABLE.LIB contains several equates pertaining to the hardware. These are explained with comments in the library file, the following is an expanded description of each item.

'DEFALT' is the drive number which will be assigned to the first volume when WINDRV is loaded, usually drive 3 will be volume A. Later, if you are going to boot FLEX from the winchester drive, the default drive must be changed to zero to allow access to STARTUP.TXT.

'DELAY' is used to time out waiting for the controller to indicate ready after indicating busy. This is only used when the controller is not powered up. If your drive is always powered when the computer is on, you can set the delay to 10000 and forget it. Too short of a delay may cause 'DRIVE NOT READY' on the first Winchester access after power up. If you power down your controller when not using your drive, the time out will prevent hanging the system if you access the drive. Repeat: most users should use 10000 and forget it! 'DRV1ST' is the number of the first (or only) physical drive. One would expect this to be drive 0, later I'll explain a small advantage to having the first drive select at 1. 'DRVORG' is the start of the drive table and the disk drivers. This location must be 'out of the way', where FLEX programs will not overwrite the drivers. If you have the interface board by Graves it has a RAM chip to hold the disk drivers. If you have an SWTP MP-09 there is RAM space on the

CPU for the drivers. Usually you will want the drivers located in the $E800 or $F000 area. WINDRV only needs about 300 bytes. 'ECFLAG' is a flag used for selecting error correction on the controllers that support it such as WD1002.05. If you have a WD1000-05, then set the flag to $00. If you have a WD1002.05, you will want error correct enabled, set the flag to $80. 'HBASE' is simply the starting address of the interface. 'HEADS' is the number of data heads per drive. For example the SA-604 has 4 heads for its two platters. 'INTERL' is the default interleave factor. You may override this number by letting WINFMT prompt you for interleave. You will need to experiment with interleave factors to find the optimum number for your system. You can set this to 16 for now. 'MAXDRV' is the largest drive number your FLEX/STARDOS system can handle. Most FLEX systems are limited to 0-3, STARDOS and my FLEX allow drives 0-9. 'NODRIV' is the phantom drive selected to turn off all the drive select LEDs. If you don't have a third drive, use 3. The driver will select this number to turn off all lights. 'OFFSET' is a trick to allow old motherboards to support eight address ports. If you have the Graves board, or you have at least 8 addresses per port, set this equate to $00. If you have 4 addresses per port (like me) then you can still use the interface by connecting A5 to an unused I/O buss line (such as UD3). This will result in the board decoding 8 addresses, though in two groups of 4. My card responds to $E010-$E013 and $E030-$E033. Setting OFFSET to $1C allows access to this split address. 'PRKTRK' is set to slightly over TRACKS. The drive manufacturers build the drive with extra tracks to provide a 'landing zone' where you can safely park the heads. Most users allow the heads to fall wherever they were at power off, usually on top of a directory! 'PRECMP' is set to 32 for most drives. This track multiplied by four is the start of precompensation. 'SEEKSP' is set to 0 for most drives. See the section SEEK and RESTORE for more information. 'TRACKS' is simply the number of tracks per drive. For example the SA-604 has 160 tracks so TRACKS is 160.

The latter section of WINTABLE.LIB contains 26 volume size equates. The numbers are the size, in DATA tracks, of the volume for each letter. If SIZA is equated to 30 then volume A will have 30 data tracks and 1 directory track (31 total). The directory of volume A will be at track zero and the data tracks will be tracks 1-30. If SIZB is equated to 70 then volume B will

have its catalog at track 31, immediately after volume A. The 70 data tracks for volume B will be physical tracks 32-101. The total number of tracks (SIZA+1,SIZB+1,etc) must add up to multiples of the number of tracks per drive. If not, WINSYS will give an error message. In other words, a volume can't extend across two drives. If you have more than one drive, and a volume ends exactly at the end of the first drive, then the next volume will start on the next physical drive. You needn't use all the tracks on a drive, WINSYS will not complain if the end of the last volume is less than TRACKS. You may skip any volume letters by equating them to a size of zero.

## WINTABLE.LIB EXAMPLES

All this should become clearer when you look at the examples. EXAMPLE A1 is from a simple system, and is the file that is distributed as WINTABLE.LIB. The first half should be easy to follow. The controller is a WDI000-05 connected to a Graves board, so DRVORG, ECFLAG, OFFSET, and HBASE are set accordingly. The drive is a 5 MEG ST-506 compatible, so TRACKS, PRKTRK, and HEADS are set accordingly. There is only one drive which is jumpered to be selected as zero, so DRVIST and NODRIV are set accordingly.

The size table is set up with only three volumes. They are selected to total 160 TRACKS. Note that the sum must include the directory tracks:

SIZA SIZB SIZC MAX 30+1 + 70+1 +57+1 = 160

EXAMPLE B1 is from a more elaborate system, mine! The hardware equates are set up for 2 drives (jumpered as 1 and 2). The driver is addressed in a small space in my version of FLEX. The error correct feature of the controller is enabled. The maximum drive number on my system is 9, so that I can have 10 Winchester volumes enabled at once (or 4 floppies, 5 Winny volumes, and a RAMDISK). All but 2 volume letters are used, note that I have added the volume names and drive numbers as comments. Sometimes I could match the letter to the name (S=SOURCE, P=PICTURE). I have duplicate volumes on both drives because my main purpose for the second drive is quick backups.

## FORMATTING DISKS

We're getting down to business! Now that you have set up the tables you can assemble WINFMT. There are no options to edit in WINFMT, it gets the user defined equates from the library file. WINFMT is self contained, you do not need to install the drivers before formatting. After assembly, enter the command WINFMT (do not enter any command line parameters for now). With luck you will see a display like EXAMPLE A2. You may get an error message on the bottom of the display if WINFMT does not like your table. There will be 3 items reported for each volume letter. The size is the same number you entered in the equate table. WINFMT has computed the drive number (always zero in EXAMPLE A2) and the starting track number for the volume. If there is an error in your table (such as a volume extending past TRACKS) the bad volume will be marked with asterisks. All volumes following the error will be marked as bad, since they are computed relative to the bad volume. WINFMT will abort if there is an error, you must fix the table and reassemble before continuing.

When the program is happy with the table, it will prompt for the volume letter to be formatted. The first time through, start with volume A. Note that any volume can be formatted without disturbing the others. The next prompt is for the number of surfaces. You should enter the number of surfaces supported by your drive, though you may enter less for testing purposes. Some drives use one surface for internal servo information, do not try to format that surface! If you enter too high a number, WINFMT will run VERY slow since it is getting errors on the non-existent head. Interleave factor is the next prompt, enter 16 for now. Interleave is discussed in detail later.

The program will now attempt to read the old volume ID, if none exists there will be a pause of a couple of seconds. If the volume has been formatted before, the old name and number will be reported. This is to reduce the chance you will format a volume you need; you can exit WINFMT by pressing return. You will now be prompted for the volume name. This has the same limitations as a FLEX file name: 8 characters, first must be alpha, hyphen and underline allowed, 3 character extension allowed. The last prompt will be for volume number which can be up to 65535.

WINFMT will now proceed with the formatting. There will be a pause as the drive is restored. Empty sectors will be created on the disk, then the sectors are filled with zeroes and the FLEX link. The program will display track numbers as it proceeds to let you know how it is progressing. The numbers should count quickly, if the tracks take over 2 seconds each there is something wrong; check the number of heads, or the head data from the 20 pin cable. Unlike a floppy, when you change the drive select jumper you must also move the head data cable to a different jack on the controller. The speed of the linking phase will vary depending on the interleave, but please do not judge the interleave factor by the format speed! The interleave should be optimized for reading files. If all goes well the volume will be formatted. Try formatting the volume again to see if the volume name and number is reported! At this point you can format the other volume letters, though you may format them several times before the day is over! You might format some of the other volumes using WINFMT with a command line parameter. To format volume C without having to answer several prompts, you can type WINFMT C. The number of heads and the interleave factor will be taken from the library file to save you some typing.

## INSTALLING THE DRIVERS

Assemble WINDRV next. The drivers file is set up as a command file, you load the drivers by calling them as a command. Later, you may want to append them to FLEX; for now you can call them manually or add WINDRV to your STARTUP file. WINDRV will load into two sections of memory. The driver itself will be located at the address DRVORG found in the library file. There will also be a one time initialization program loaded at $C100 which does its thing then jumps to the FLEX cold start. The program will save the old FLEX driver jump table (usually the jumps to the floppy routines) then install its own jump table. Unlike the Zeff-Graves program you do not need to change your winchester driver when you change your floppy driver. Also, you can load another driver after this one such as my RAMDISK.

If all is going well you should be able to catalog the volume A that you formatted earlier (using CAT 3). WINDRV assigns the first letter to drive 3, since there will always be a volume

starting at track zero. Unless you equated SIZA to zero, CAT 3 should display a catalog of volume A. It is assumed you will want to change this drive 3 = volume A pair with WINASN after the driver is loaded. I included this default assignment for test purposes so that you can test the drivers without WINASN.

## ASSIGNING WINCHESTER VOLUMES

The next program to assemble is WINASN. Again, no changes need to be made to this file. This program is used to select a volume letter to be used with a drive number. After loading the drivers you could only access volume A, which was assigned to drive 3. Enter the command WINASN without any parameters. You should get a list like EXAMPLE A3. This varies from the WINFMT report in that it displays the volume name (and extension if you entered one). This display tells you what volumes are available, sort of a directory of directories. I won't go so far as to say "UNIX like", but a small step in that direction. Along the bottom of the display is a list of drive numbers with their assigned volume letters. 0=- means no volume assigned, 3=A means volume A is assigned to drive 3. WINASN prompts for drive number/volume letter pairs. If you want to assign drive 1 to volume B enter 1B as an answer to the prompt. Now if you enter CAT 1 you will catalog winchester volume B.

Several questions may be popping up in your mind... What happened to my floppy disk at drive 1? The floppy temporarily vanishes from access by FLEX. All references to drive 1 will access the winchester volume assigned to drive 1. If you want the floppy back, you can enter WINASN 1-. Remember I mentioned 0=- is WINASN's way of indicating no volume assigned? The dash (hyphen) is also used to deassign a volume. Does that mean I can assign all my drive numbers as winchesters? Sure does! Try the command WINASN 1A 2B 3C. Assuming you formatted the first three volume letters, you now have three

winchester volumes assigned (for now you don't want to assign a winchester drive as the system drive). Now try CAT 1 2 3 and see what appears. What does WINASN accept as parameters? The program wants (and insists) on number/letter pairs. There can be several pairs on one command line. Commas and spaces can be inserted if you find that more readable. For example WINASN 0 A 2,B 3C,4D5E is perfectly acceptable. Since I have 4 floppies my STARTUP file includes WINASN 4C5G6L7H8W3-, the latter two characters re-enable my floppy at drive 3. What other special characters does WINASN have? Like WINFMT the asterisk (star) is used to mark a volume whose size must be corrected. A pound sign (pig pen) is used to indicate a drive not responding when WINASN tries to read the volume ID. How can I remember the special characters? By remembering another special character of course! I think you all know the familiar plus sign, if used as a parameter you will get a help message. Enter WINASN + to get some help on how to use this command. Must the letters be in any order? No, WINASN 1C,0B,2A,3C will work fine though there is little purpose to use the letter C twice. How does this relate to the FLEX ASN command? The two perform separate functions, but work together as a team. WINASN 3C:ASN W=3 will make volume C the work drive. How long will it take to read all those volume names? I find it takes about 2 seconds for WINASN to list my 24 volumes scattered around two 5 megabyte drives. Until all volumes have been formatted, the command will take longer.

Since this is the command you will use the most, I'd suggest playing with it for awhile. It's really very intuitive since you simply tell it the drive number followed by the volume you want to assigned to that drive.

## PARKING A WINCHESTER

Most people simply turn off their computers when done with them. Unfortunately winchesters do not have any hardware that holds the flying head away from the platter. The head just lowers down as the disk slows, finally coming to a scraping stop! Thousands of disk users do this every day, and little is said about the subject. In my profession of servicing computers I've changed my share of winny drive mechanisms, often as the customer mumbles about how long it's been since the disk was backed up! I don't know if any of the failures were due to not parking, but I'd rather not find out the hard way. When a winny fails, it usually goes all at once without any warning. The makers of the drives have always provided landing zones, a section of the disk that is not used for data. They recommend resting the heads there, especially if the drive is to be moved. I've noticed IBM is now providing parking support for their PC-AT, H-P does likewise for their VECTRA. Perhaps PC-XT users have learned the hard way...

WINPARK is the final program for you to assemble. You may use it if you like, ignore it if you must. There is one option in the source: a flag to enable reporting of the drive and track being parked. I use the command before powering the system down.

*FOR THOSE WHO* NEED TO KNOW **68 MICRO JOURNAL™**

## MICROWARE UPDATE

The following information was extracted from various Microware publications and are presented for your information.

### OS-9 VERSION 2.2

Microware has released a new 2.2 Version of the Industrial, Personal and Professional OS-9 Operating System. The new version includes an update of selected operating system modules, plus an enhanced version of Microware's C Compiler. This release does not require the re-assembly or modification of any system-level software modules or application programs. Most OEM's have completed integration and testing of V2.2 and are currently shipping this new version of OS-9.

The module updates includes additional functionality to support three exciting new Microware software products (previewed in the Fall 1987 edition of Pipelines): C Source Level Debugger, Ethernet TCP/IP package and Memory Protection Module (SPU). Of course, the release contains corrections for errors reported in the 2.1 Version of OS-9/680x0.

New features added to C include bit-fields, enumerated data types and the ability to return structures and unions. The use of bit fields can reduce program memory requirements by allowing data to be stored in a more efficient sized variable. Enumerated data types have also been included. Their use makes programs more understandable and error free by permitting constants to be referenced by a symbolic name. And finally, the ability to return structures and unions is useful for manipulation of variables. With the addition of these three new features, Microware's Version 3.0 C Compiler now represents an even more powerful C Language tool for software development.

All changes included in the new OS-9 Version 2.2 are documented fully in Release Notes that accompany the update disks. End users may refer to these notes for a more exact description of changes made in the operating system. Anyone having questions regarding V2.2 should contact their Microware representative or their system supplier for additional information.

## NEW DEVELOPMENT PAKS FROM MICROWARE MVME 133A, 134 &135 NOW AVAILABLE

New Development Paks for Motorola VME 133A, 134 and 135 CPU's are now available from Microware. The MVME 133A, 134 and 135 are popular 68020 CPU's designed for software development and industrial applications. The new Development Paks have been customized to support both European (MVME 319) and American (MVME 320) configurations. An optional MVME 050 system controller can be added to provide two additional serial ports and one parallel port in addition to other features.

The MVME 133A includes complete device driver or trap handler support for the on-board MC68881 FPU, RS-232 serial ports and 20 Mhz clock. These features make this micro-computer a popular choice for numerically intensive applications. The MVME 134 has a MC68851 Paged Memory Management Unit supported by Microware's OS-9/SSM Security Pak making it an ideal engine for large multi-user systems. The significant feature of the MVME 135 is a VSB private memory bus which OS-9 supports for multi-processor systems and systems incorporating high-speed DMA controllers. As with all Microware/Motorola Development Paks, these versions of OS-9 can support up to 16 megabytes of RAM. Microware has also recently released an Optional MVME Driver package for all Motorola VME Development Paks. The Optional Driver package includes object code for the MVME 335 serial/parallel device drivers and the MVME 350 magnetic tape device drivers.

## MICROWARE WELCOMES ITALIAN AUTHORIZED IMPLEMENTORS

Microware and Eximar S.r.l. of Rome, Italy, have signed an Authorized Implementors agreement for OS-9 technical sales and service in Italy. As Authorized Implementors, Eximar will be working with Microware products and its Distributors to promote Microware products throughout Italy. Their primary focus will be to provide a complete menu of technical assistance including education, software support, integration and installation services.

Eximar's resident, Carlo Narcisi, previously served in the Italian armed forces and represented Motorola GEG for 15 years prior to forming Eximar in 1983. Eximar's Technical Services manager, Ercole Maccioni, also worked for Motorola Microsystems prior to joining the firm in 1984. Ercole and Roberto Evangelista came to Microware for technical training and product marketing meetings the week of April 18th.

Eximar services the Italian defense electronics industry on behalf of selected American high-technology companies. They also provide installation and integration services for commercial customers in the process control and board-level markets.

Readers interested in more information may contact Eximar at the following address:

*Eximar S.r.l.*
*Via Dei Ciceri, 59 • 00175 -Rome, Italy*
*Phone: (6) 762221 • Telex: 623358*

## ADA CROSS COMPILER NOW AVAILABLE FOR OS-9

Microware announces the availability of the Systeam Ada Cross Compiler for VAX/VMS systems from our West German distributor, Dr. Rudolf Keil, Gmbll. The Systeam Ada conforms to the Department of Defense 4.4 Ada specification. Microware now can offer its customers a tool to pursue U. S. government contracts which specify software projects to be developed under the Ada programming language.

The Ada Cross Compiler is hosted on VAX/VMS and produces object code for a Motorola MC68020 CPU and MC68881 FPU running under OS-9. It translates full Ada as specified in the 1983 Ada Reference Manual.

The user interface is easy to use. The cross compiler is well integrated into the underlying host operating system. It can recognize about 500 different English language error messages. Messages are self-explanatory and accurately positioned.

The cross compiler can translate up to 1800 lines per minute on a VAX 8500 and requires 1.5 MBytes of hard disk storage space. The amount of virtual memory required for a given application depends on the size and number of compilation units to be compiled and typically varies between 4 and 7 MBytes.

The cross compiler is delivered on magnetic tape or floppy disk and full documentation is supplied with every cross compiler. To order the Systeam Ada Cross Compiler, or for more information, contact Dr. Keil at the address listed below:

*Dr. Rudolf Keil, GmbH*
*Gerhart-Hauptmann-Str. 30*
*6915 Dossenheim, West Germany*
*Phone: 06221-862091*
*Telex: 461166*
*FAX: 8221-861954*

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## ATARI & AMIGA CALL

As most of you know, we are very sensitive to your wishes, as concerns the contents of these pages. One of the things that many of you have repeatedly written or called about is coverage for the **Atari & Amiga™** series of 68000 computers.

Actually we haven't been too keen on those systems due to a lack of serious software. They were mainly expensive "game-toy" systems. However, recently we are seeing more and more honest-to-goodness serious software for the Atari & Amiga machines. That makes a difference. I feel that we are ready to start some serious looking into a section for the Atari & Amiga computers. Especially so since OS-9 is now running on the Atari (review copy on the way for evaluation and report to you) and rumored for the Amiga. Many of you are doing all kinds of interesting things on these systems. By sharing we all benefit.

This I must stress - *Input from you on the Atari & Amiga*. As most of you are aware, we are a "contributor supported" magazine. That means that YOU have to do your part. Which is the way it has been for over 10 years. We need articles, technical, reviews of hardware and software, programming (all languages) and the many other facets of support that we have pursued for these many years. Also I will need several to volunteer to do regular columns on the Atari & Amiga systems. Without constant input we can't make it fly! So, if you do your part, we certainly will do ours. How about it, drop me a line or give me a phone call and I will get additional information right back to you. We need your input and support if this is to succeed!

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